Networking Technologies and Applications

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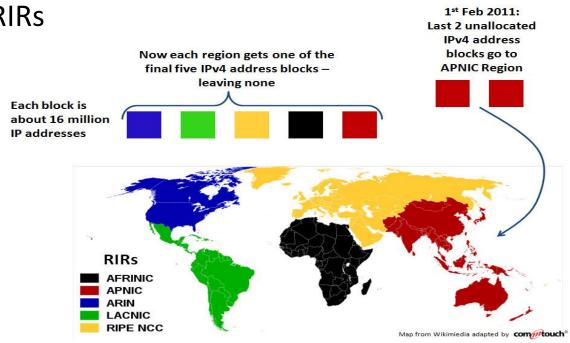


Exhaustion of IPv4 addresses

- No problem in the US
 - "Internet Heaven"
- Serious problem everywhere else
 - In China they have asked IP addresses for connecting 60.000 schools, they received a Class B address (65.534 addresses)
 - Many European or African countries received just a Class C address (254 addresses)
- Fast development of the Internet outside Northern America
 - Asia (2.5 billion people), Eastern Europe (250 million), Africa (800 million), South and Latin America (500 million)
- New communication devices need IP addresses
 - Mobile phones, PDAs, sensors, cars, etc.
- The exhaustion of IPv4 addresses was always projected for the next month/years (for more than 10 years)

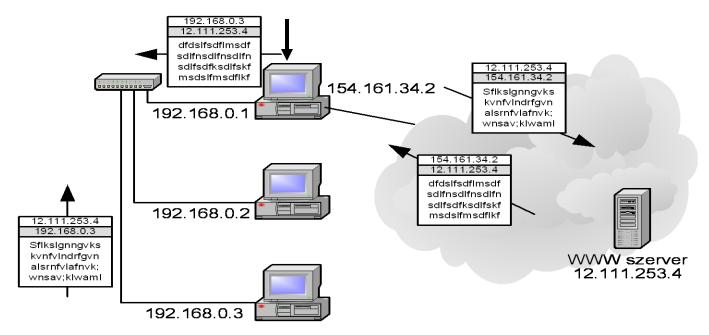
Is the Internet "full"?

- On February 1st, 2011, the last /8 IPv4 address block allocated to the RIRs
- Last allocated blocks by RIRs
 - APNIC April 2011
 - RIPE Sept. 2012
 - LACNIC June 2014
 - ARIN Sept. 2015
 - AfriNIC 2018?



Reuse of private addresses

- Private addresses are not "visible" from the Internet
- Network address translation



NAT problems

- Just an intermediate solution
 - Cannot establish a connection from outside with a host behind a NAT
 - More and more applications require public IP addresses
 - VoIP, videoconferencing, network games
 - Many protocols do not work on a network with NAT

IPv6 chronology

- TUBA (1992)
 - TCP and UDP over Bigger Addresses
 - Based on the OSI CLNP (Connection-Less Network Protocol)
 - abandoned
- **SIPP** (1993)
 - Simple IP Plus
 - 64 bit addresses
- IPng, based on an extended SIPP version (1994)
 - 128 bit addresses
 - From December 1995 officially called IPv6

IPv6 addressing scheme

- The IPv6 address pool is huge
 - $-2^{128} = 340.282.366.920.938.463.463.374.607.431.768.211.456$
 - 67 billion billion addresses for each cm² on the surface of the Earth
 - 10³⁰ address for each person on the Earth
 - The address distribution and routing requires a hierarchical structure

IPv6 addressing scheme

- IPv6 addressing quite similar to IPv4
 - 128 bit long addresses, instead of 32 bits
- Three address types:
 - Unicast addresses
 - Identify a unique interface
 - Multicast addresses
 - Identify a group of interfaces, each of these will receive the message
 - Replace the broadcast addresses as well
 - Anycast addresses
 - Identify a group of interfaces, message will be delivered to one of these interfaces

Writing IPv6 addresses

- 128 bits = 16 bytes = 32 x 4bits = 32 hexadecimal digits grouped in 8 segments FECD:BA98:0000:0000:00CD:BA98:0000:3200
- The opening 0 digits in each segment can be neglected
 Instead of FECD:BA98:0000:0000:00CD:BA98:0000:3200

we write FECD:BA98:0:0:CD:BA98:0:3200

- Adjacent 0 segments can be neglected, if there is only one siuch case in an address FECD:BA98::CD:BA98:0:3200
- Network prefix is encoded as in case of IPv4 CIDR

entire IPv6 address/prefix length in bits

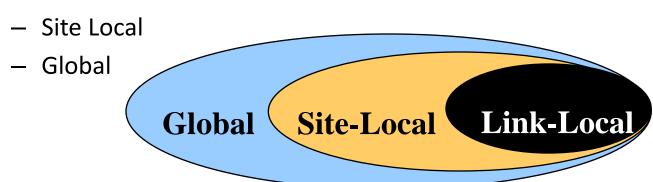
12AB:0000:0000:CD30:FFFF:DEC8:0000:0000/60

12AB:0:0:CD30:0:0:0:0/60

12AB:0:0:CD30::/60

IPv6 addressing scheme

- One interface might have several addresses, with different scopes:
 - Link Local



Unicast addesses

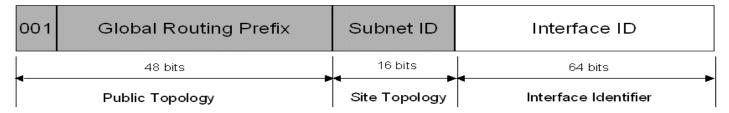
- Valid over a limited scope
 - Scoped address
 - Novelty in IPv6
- Scope = local link
 - For communication among nodes on the same link
 - Unique only on that link, cannot be used for communication outside the link
 - Automatically configured on each interface
 - Each IPv6 node has an initial address with which it can start communicating
 - Neighbor dicovery, router discovery
 - Format:
 - FE80:0:0:0:<interface identifier>
 - Interface ID EUI (64) address
 - Extending the previous 48 bit MAC address

Unicast addresses

- Scope = site local
 - Used for communication inside the same site
 - Routers will not forward it outside the site (to the Internet)
 - Similar to IPv4 private addresses
 - Not automatically configured
 - Format:
 - FEC0:0:0:<subnet id>:<interface id>
 - Subnet id = 16 bit = 64K subnet
 - Allows the addessing of an entire organization (company, university)
 - E.g. allocate site-local addesses to the entire network
 - Renumbering when connecting to an IPv6 network
 - We change the first 48 bits (to a site ID)
 - Renumbering when connecting to a new service provider

Unicast addresses

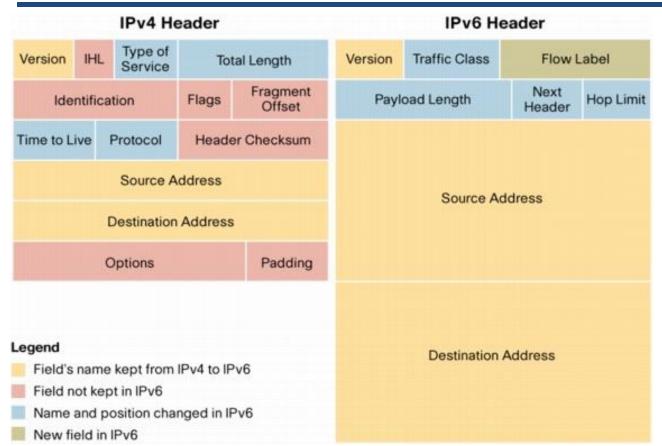
- Global Unicast Address
 - Used for global communication
 - Hierarchical global prefix
 - Structured by RIRs and ISPs
 - Subnet ID
 - Hierarchically structured by the network administrator
 - Interface ID



Multicast addresses

- Instead of broadcast addresses
- Scoped addresses
 - Node, link, site, organisation, global
- Format
 - FF<flags><scope>::<multicast group>
- Flag:
 - 0 permanent
 - 1 dynamic
- Scope:
 - 1 node
 - 2 link
 - 5 site
 - 8 organisation
 - E global
- E.g.
 - FF02::1 all nodes on the local network
 - FF02::2 all routers on the local network

Format of the mandatory fixed IPv6 header



IPv6 mandatory fixed header

- Version (4 bits)
 - IP version
- Class priority class (8 bits)
 - Defines the priority of the packet
 - DSCP field in IPv4
 - Used for providing Quality of Service (discussed later)

IPv6 mandatory fixed header

Flow Label

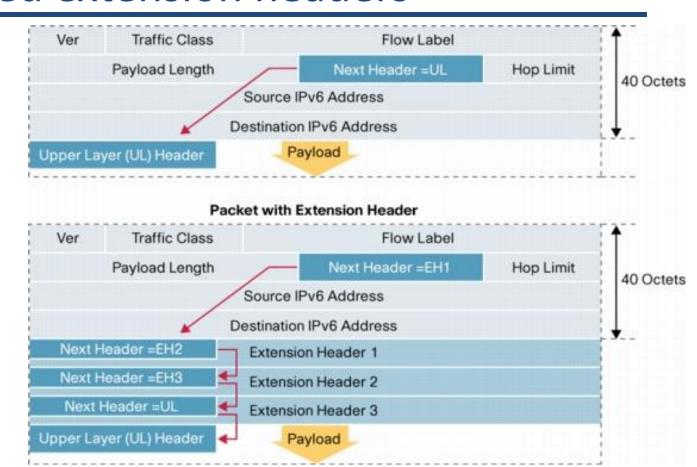
- Associated to traffic with special QoS requirements
 - 20 bit long
- Can be used in router caches to decrease the processing time
 - A packet first arrives to the router
 - The fow label is stored in the cache
 - When the next packet with the same flow label arrives...
 - No need to look up the destination address in the routing table
 - Packet immediately routed based on the flow label
- Payload Length (16 bits)
 - Length of the useful data, in bytes

IPv6 mandatory fixed header

- Next-Header (8 bits)
 - Identifies the header directly following the fixed header
 - It might be an extension header, or a higher layer protocol header (e.g., TCP)
- Hop Limit (8 bits)
 - Provides how far a packet can travel
 - Same as the Time To Live (TTL) field in IPv4
- Source Address (128 bits)
 - Address of the original source of the packet
- Destination Address (128 bits)
 - Not necessarily the address of the last receiver, if the packet also contains a Routing Header

- IPv6 packet starts with a 40 byte long mandatory fixed header
- Extra information related to the intermediate network encoded in Extension Headers
- Most of the extension headers are not processed by the intermediate routers, only by the final destination
- Each extension header has a special value for the next header field
 - Many extension headers can be used in parallel
 - The value of the last next header field identifies the upper layer protocol
 - The header can be of any length

IPv6 chained extension headers

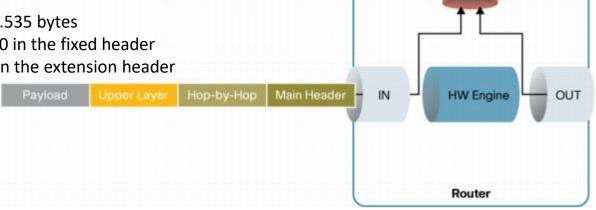


Header order

- Important: extension headers should respect the suggested order
 - Easier for routers to process the packet
 - In most cases the routers process only the hop-by-hop options and the routing header
- Exception: destination option
 - Right before the higher layer header
 - If we want the intermediate routers to process the destination option header, we should put it right before the routing header, and they should be processed together.
 - A packet might contain a destination option headers in both locations

- The suggested header order:
 - IPv6 Header
 - Hop-by-hop Options Header (type = 0)
 - Destination Options Header (1)
 - Routing Header (type = 43)
 - Fragment Header (type = 44)
 - Authentication Header (type = 51)
 - Encapsulating Security Payload (ESP) (type = 50)
 - Destination Options Header (2) (type = 60)
 - Upper Layer Header (pl. TCP vagy UDP)

- Hop-by-hop Options Header
 - Contains IP options for the intermediate routers
 - Each intermediate router should analyze and process the Hop-by-hop Header
 - Router Alert option alerts tranzit routers
 - If the packet packet contains information that should be processed by an intermediate router
 - Otherwise the packet is not analyzed, just routed
 - IPv6 jumbogram option
 - For packets larger than 65.535 bytes
 - The payload length set to 0 in the fixed header
 - The true length specified in the extension header



Process the

Hop-by-Hop EH

CPU

Routing Header

- In the normal case the source of the IP packet leaves the routing task to the network
- In case of source routing, the source will specify a path with router addresses
 - The entire list in the Routing Header (e.g., A, B, C, D)
 - The destnation address is always the address of the next router specified in this field, except the last router
 - Each router modifies the destination address, before forwarding the packet

- Fragment Header
 - In IPv4 fragmentation and reassembly automatically, if explicitely not forbidden
 - Don't Fragment flag
 - In IPv6 packets are not fragmented by default
 - If the packet is too large for the transmission medium, it is dropped and an ICMP (Internet Control Message Protocol) error message is sent
 - The source discovers the path MTU-t
 - Maximum Transmission Unit
 - Tries to send packets with a lower size than the MTU
 - If we need fragmentation, that can be done using the Fragmentation Extension Header
- Authentication Header
 - Guarantees that ...
 - The recieved packet is authentic
 - It was not altered on its path
 - Comes from the specified source
- Destination Option Header
 - Contains options to be processed by the destination node