

Networking Technologies and Applications

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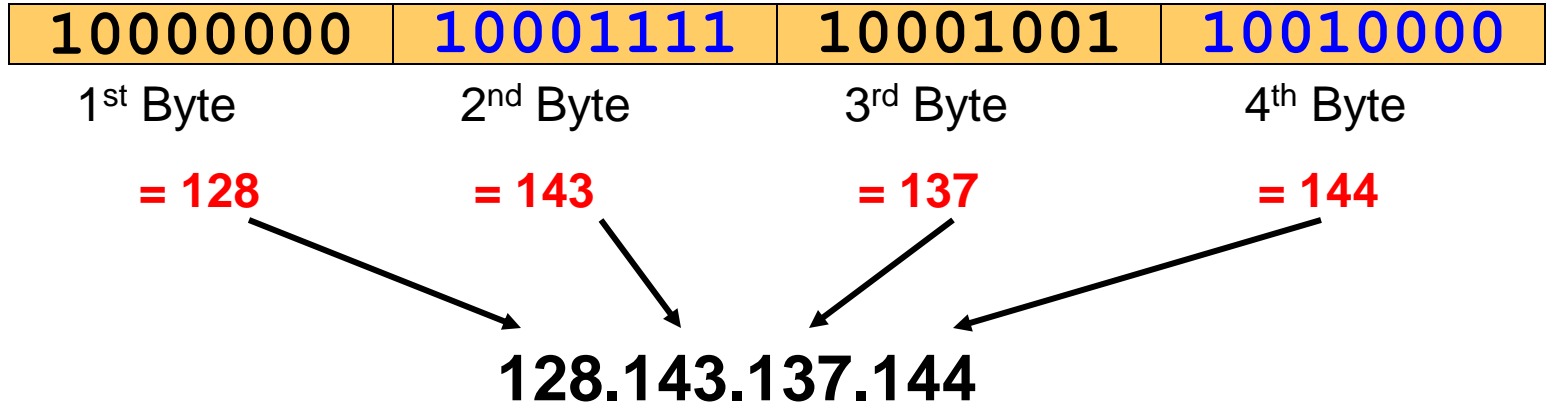


What is an IP Address?

- An IP address is a unique global address for a network interface
- An IP(v4) address:
 - is a **32 bit long** identifier
 - encodes a network number (**network prefix**) and a **host number**

Dotted Decimal Notation

- IP addresses are written in a so-called *dotted decimal notation*
- Each byte is identified by a decimal number in the range [0..255]:
- **Example:**



Network prefix and Host number

- The network prefix identifies a network and the host number identifies a specific host (actually, interface on the network).

network prefix

host number

- **How do we know how long the network prefix is?**
 - The network prefix used to be implicitly defined (**class-based addressing, A,B,C,D...**)
 - The network prefix now is flexible and is indicated by a **prefix/netmask (classless)**.

Example

Example: argon.cs.virginia.edu

• IP address is 128.143.137.144

- Is that enough info to route datagram??? -> No, need netmask or prefix at every IP device (host and router)

• Using Prefix notation IP address is: **128.143.137.144/16**

- Network prefix is 16 bits long

• Network mask is: 255.255.0.0 or hex format: **ffff0000**

-----> **Network id** (IP address **AND** Netmask) is: 128.143.0.0

-----> **Host number** (IP address **AND** inverse of Netmask) is: 137.144

128.143

137.144

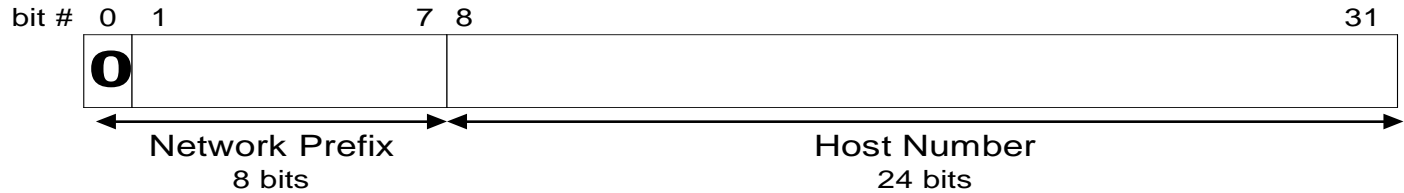
The old way: Classful IP Addresses

- When Internet addresses were standardized (early 1980s), the Internet address space was divided up into classes:
 - **Class A:** Network prefix is 8 bits long
 - **Class B:** Network prefix is 16 bits long
 - **Class C:** Network prefix is 24 bits long
- Each IP address contained a key which identifies the class:
 - **Class A:** IP address starts with “0”
 - **Class B:** IP address starts with “10”
 - **Class C:** IP address starts with “110”

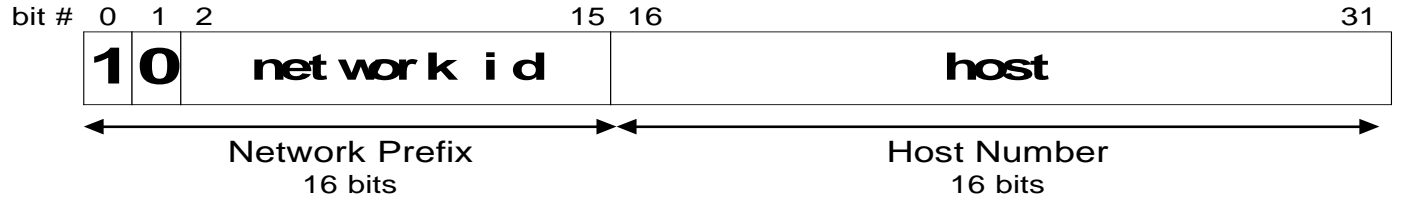
	Number of networks	Maximum nr. of hosts on a network	Value of first byte
Class A	126	16,777,214	1 – 126
Class B	16,384	65,534	128 – 191
Class C	2,097,152	254	192 - 223

The old way: Internet Address Classes

Class A



Class B



Class C



Addressing rules

- The Network ID cannot be 127
 - Reserved for the loop-back interface
- The host ID cannot be 255
 - 255 a broadcast address
- The host ID cannot be 0
 - 0 means „this network”
- The host ID has to be unique on the given network

Problems with Classful IP Addresses

- The original classful address scheme had a number of problems

Problem 1. Too few network addresses for large networks

- Class A and Class B addresses are gone
- Initially given to institutions
 - Upper left corner
 - HP, Apple, MIT, IBM, Ford, etc
- Later RIRs are created
 - Regional Internet Registrar



THIS CHART SHOWS THE IP ADDRESS SPACE ON A PLANE USING A FRACTAL MAPPING WHICH PRESERVES GROUPING--ANY CONSECUTIVE STRING OF IP's WILL TRANSLATE TO A SINGLE COMPACT, CONTIGUOUS REGION ON THE MAP. EACH OF THE 256 NUMBERED BLOCKS REPRESENTS ONE /8 SUBNET (CONTAINING ALL IP's THAT START WITH THAT NUMBER). THE UPPER LEFT SECTION SHOWS THE BLOCKS SOLD DIRECTLY TO CORPORATIONS AND GOVERNMENTS IN THE 1990's BEFORE THE RIRs TOOK OVER ALLOCATION.

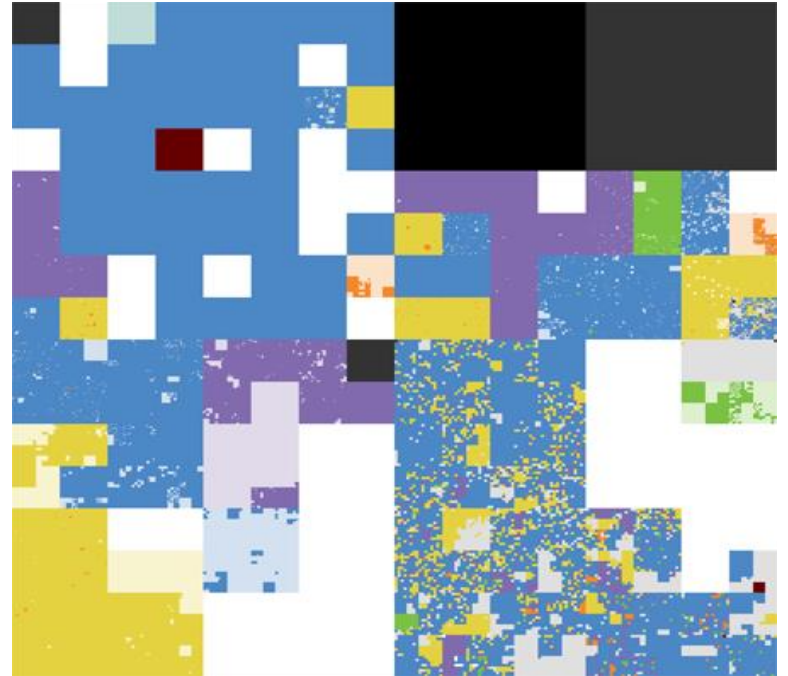
0 1 14 15 16 19 →
3 2 13 12 17 18
4 7 8 11
5 6 9 10



 = UNALLOCATED BLOCK

IPv4 addresses (2006)

- Blue: ARIN – North America
- Yellow: RIPE NCC – Europe
- Magenta: APNIC – Asia-Pacific
- Green: LACNIC – Latin-America
- Orange: AfriNIC – Africa
- Black: Multicast
- Grey: Special addresses
 - Loopback, private, class E, etc.
- White: free



Problems with Classful IP Addresses

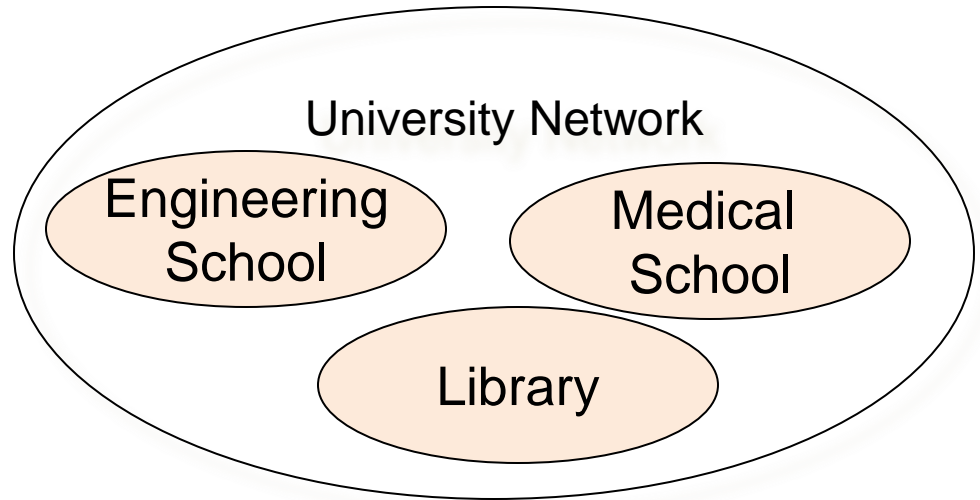
- The original classful address scheme had a number of problems

Problem 2. Two-layer hierarchy is not appropriate for large networks with Class A and Class B addresses.

– **Fix #1: Subnetting**

Subnetting

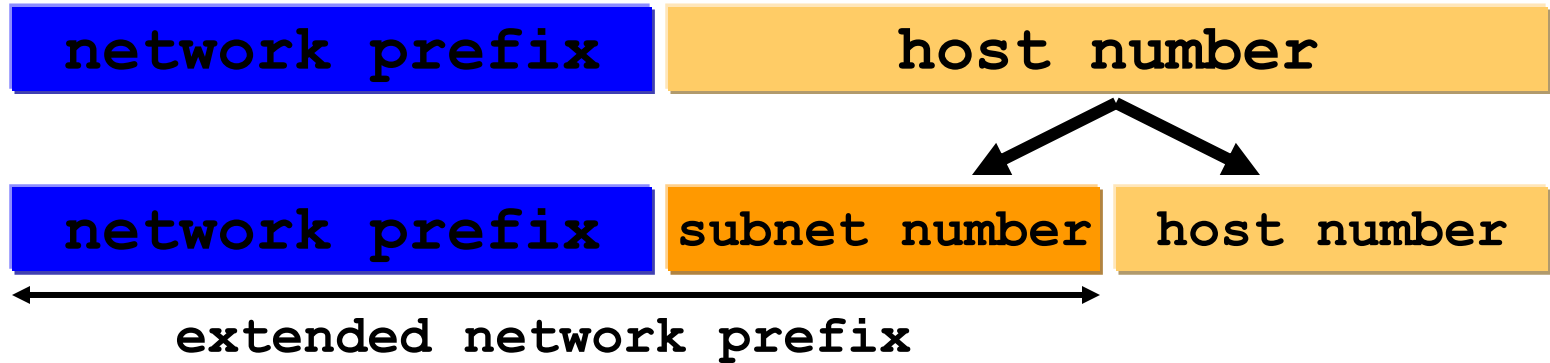
- **Problem:** Organizations have multiple networks which are independently managed
 - **Solution 1:** Allocate an **address for each** network
 - Difficult to manage
 - From the outside of the organization, each network must be addressable, must have an identifiable address.
 - **Solution 2:** Add another level of hierarchy to the IP addressing structure



→ Subnetting

Basic Idea of Subnetting

- Split the host number portion of an IP address into a **subnet number** and a (smaller) **host number**.
- Result is a 3-layer hierarchy



- **Then:**

- Subnets can be freely assigned within the organization
- Internally, subnets are treated as separate networks
- Subnet structure is not visible outside the organization

Advantages of Subnetting

- With subnetting, IP addresses use a 3-layer hierarchy:
 - Network
 - Subnet
 - Host
- Improves efficiency of IP addresses by not consuming an entire address space for each physical network.
- Reduces router complexity. Since external routers do not know about subnetting, the complexity of routing tables at external routers is reduced.
- Note: Length of the subnet mask need not be identical at all subnetworks.

Subnetting Example

- An organization with 4 departments has the following IP address space: 10.2.22.0/23. As the systems manager, you are required to create subnets to accommodate the IT needs of 4 departments. The subnets have to support 200, 61, 55, and 41 hosts respectively. What are the 4 **subnet network** numbers?
- Solution:
 - 10.2.22.0/24 (256 addresses > 200)
 - 10.2.23.0/26 (64 addresses > 61)
 - 10.2.23.64/26 (64 addresses > 55)
 - 10.2.23.128/26 (64 addresses > 41)

Problems with Classful IP Addresses

Problem 3. Inflexible. Assume a company requires 2,000 addresses

- Class A and B addresses are overkill
- Class C address is insufficient (requires 8 Class C addresses)

Problem 4: Exploding Routing Tables: Routing on the backbone Internet needs to have an entry for each network address. In 1993, the size of the routing tables started to outgrow the capacity of routers.

Fix #2 (to both of these problems): Classless Interdomain Routing (CIDR)

CIDR - Classless Interdomain Routing

- **Goals:**

- Restructure IP address assignments to increase efficiency
- Hierarchical routing aggregation to minimize route table entries

Key Concept: The length of the network id (prefix) in IP addresses is **arbitrary/flexible** and is defined by the network hierarchy.

- **Consequence:**

- Routers use the IP address **and** the length of the prefix for forwarding.
- All advertised IP addresses must include a prefix

CIDR Example

- CIDR notation of a network address:
 - **192.0.2.0/18**
 - "18" says that the first 18 bits are the network part of the address
- The network part is called the network **prefix**
- Example:
 - Assume that a site requires an IP network domain that can support 1000 IP host addresses
 - With CIDR, the network is assigned a continuous block of $1024 = 2^{10}$ (>1000) addresses with a $32-10 = 22$ -bit long prefix

CIDR: Prefix Size vs. Host Space

CIDR Block Prefix

of Host Addresses

/27	32 hosts
/26	64 hosts
/25	128 hosts
/24	256 hosts
/23	512 hosts
/22	1,024 hosts
/21	2,048 hosts
/20	4,096 hosts
/19	8,192 hosts
/18	16,384 hosts
/17	32,768 hosts
/16	65,536 hosts
/15	131,072 hosts
/14	262,144 hosts
/13	524,288 hosts

CIDR and Address assignments

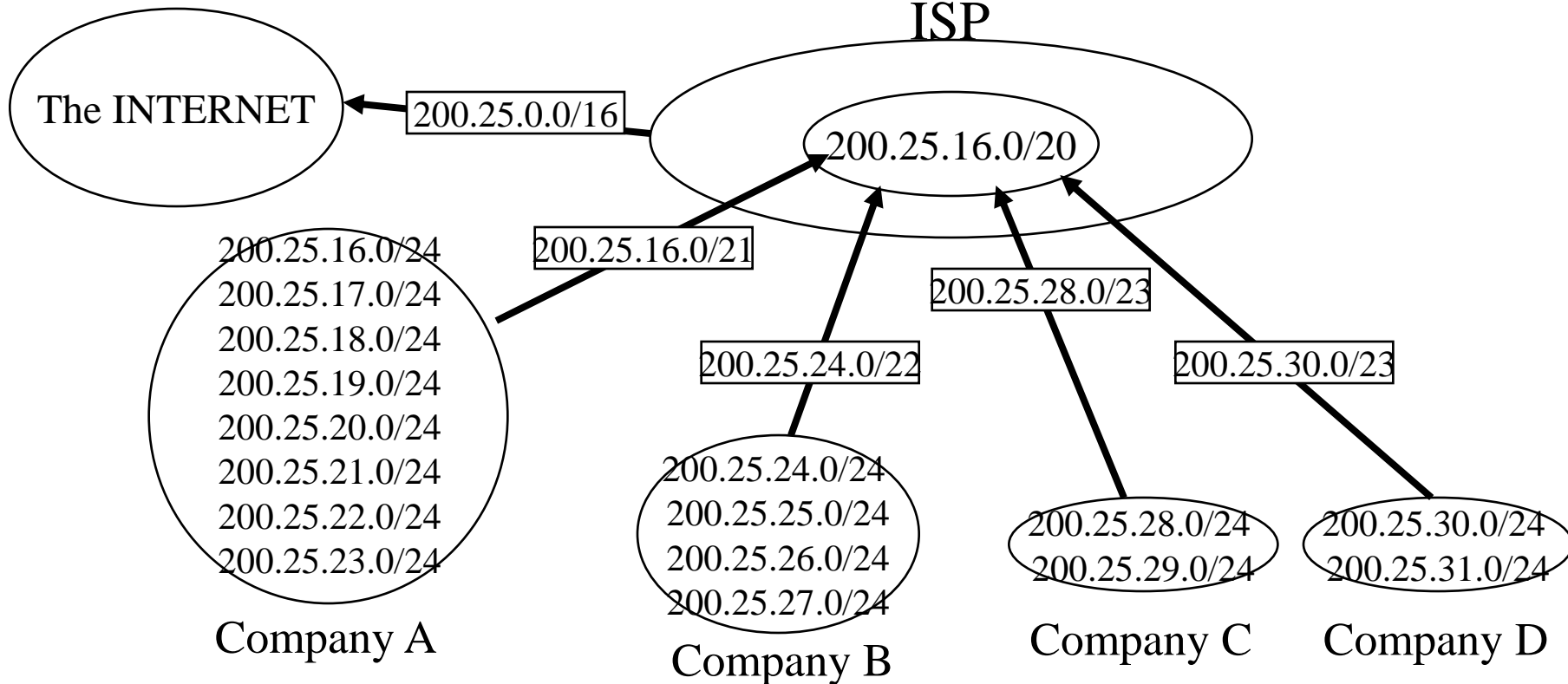
- **IANA – Internet Assigned Numbers Authority**
 - The RIRs get short prefix CIDR blocks
 - Regional Internet Registries
 - E.g., 62.0.0.0/8 assigned to RIPE NCC
 - **Réseaux IP Européens Network Coordination Centre**
- RIRs fragment and redistribute parts of the address space
 - Backbone ISPs obtain large blocks of IP address space and then reallocate portions of their address blocks to their customers.

Example:

- Assume that an ISP owns the address block **206.0.64.0/18**, which represents 16,384 ($2^{32-18}=2^{14}$) IP host addresses
- Suppose a client requires 800 host addresses
 - $512=2^9 < 800 < 1024=2^{10} \rightarrow 32-10 = 22$,
 - Assigning a /22 block, i.e., 206.0.68.0/22 -> gives a block of 1,024 (2^{10}) IP addresses to client.

01000100

CIDR example



CIDR and Routing

- **Aggregation of routing table entries:**
 - 128.143.0.0/16 and 128.142.0.0/16 can be represented as 128.142.0.0/15 at a router.
 - 143 = 128.10001111.0.0 142 = 128.10001110.0.0
- **Longest prefix match:** Routing table lookup finds the routing entry that matches the longest prefix
 - Why?

E.g., What is the outgoing interface for destination IP address: 128.143.137.0?

Prefix	Interface/outgoing link
128.143.128.0/17	interface #1
128.128.0.0/9	interface #2
128.0.0.0/4	interface #5

Routing table

Problems with Classful IP Addresses

Problem 5. The Internet is going to outgrow the 32-bit addresses

- **Fix #3: IP Version 6**