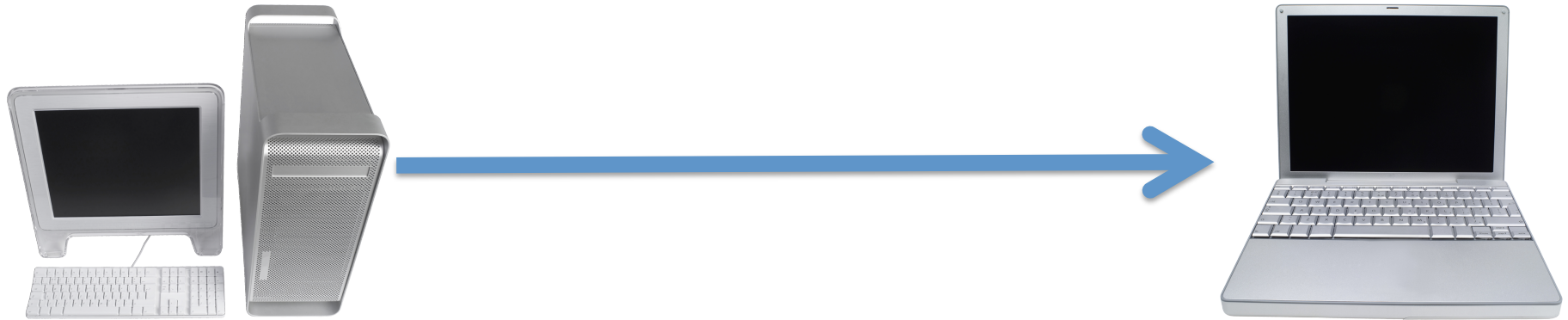


Architecture, interaction design

Agile roles

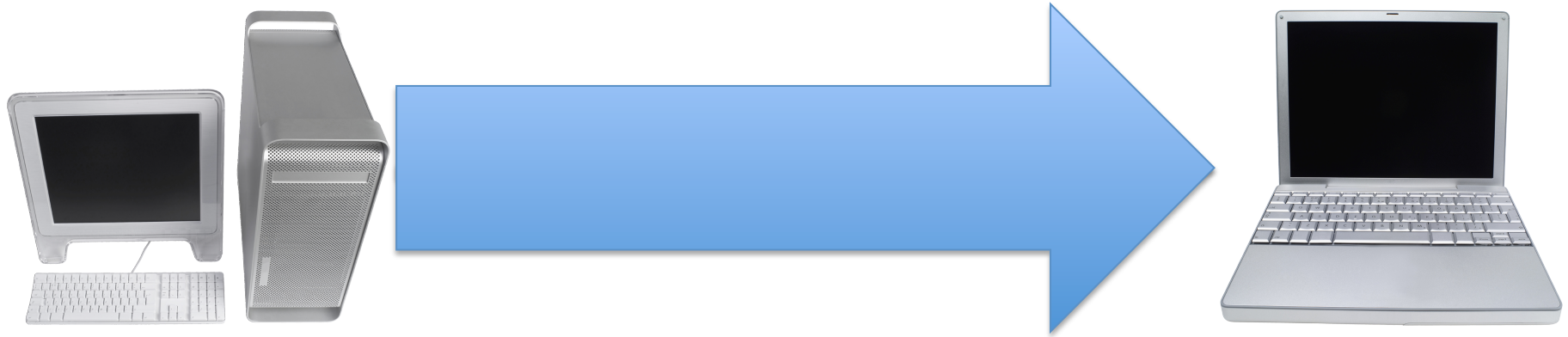
- Architect
 - A team member (programmer) who designs interfaces and data structures
- Interaction (protocol) designer
 - A team member (programmer) who designs message flows over the internal and external interfaces

What we want to do



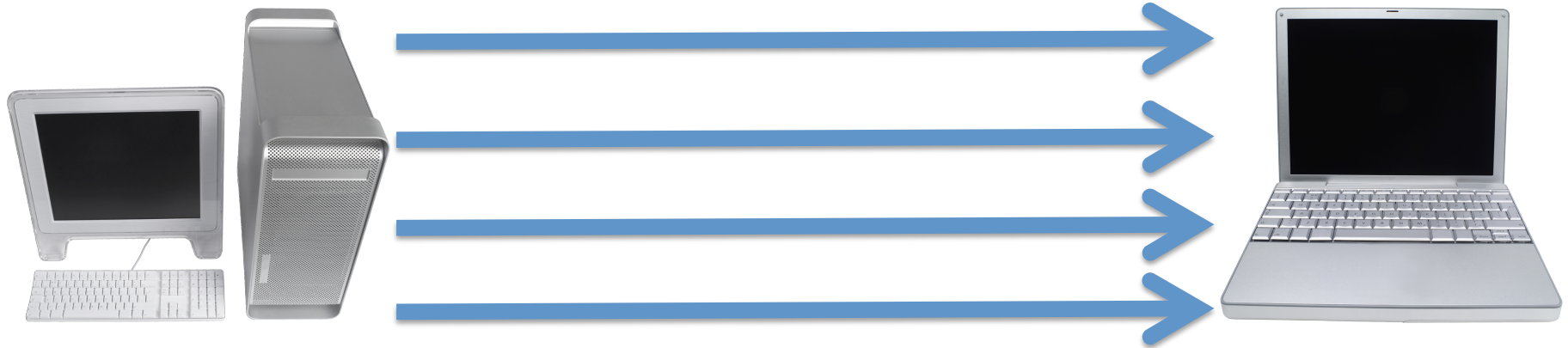
- Get data from one node to another
- In a network of a set of nodes/computers/
software component
- Taking requirements into account

All at once



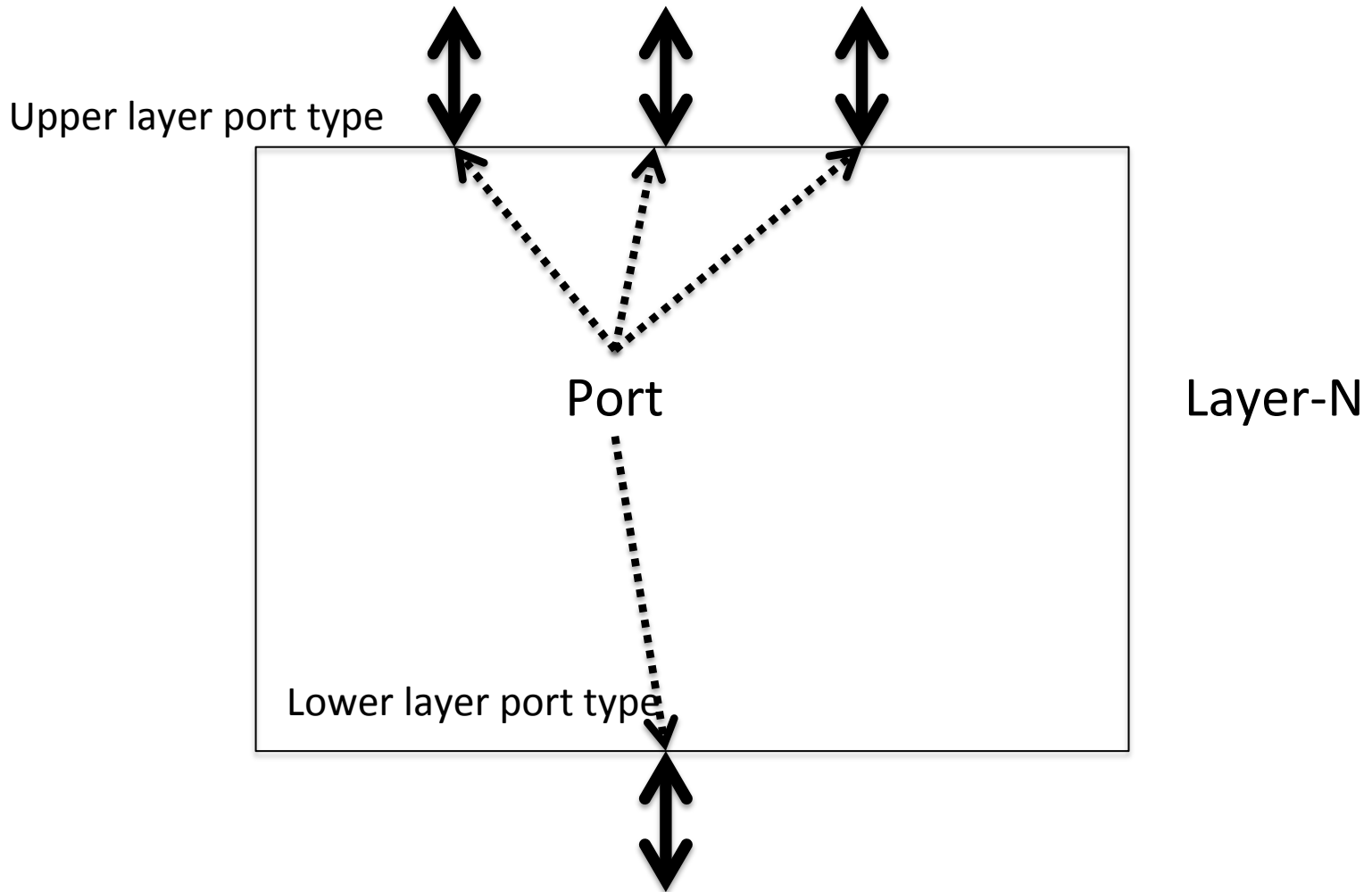
- Sending all pieces of data in one message
- Complex data structure
- Simple protocol
- Used in software, computer protocols and higher layer network protocols

In chunks



- Transfer each data item in a separate message
- Simple data structure
- Complex protocol
- Used in telecommunications networks

Network node model



Network node model

- Analogy between network node and software component
 - Layer vs. tier
 - Port type vs. interface
 - Message vs. function declaration
 - Message format vs. formal parameter list
 - Ingress port activation vs. function called or return from a call
 - Egress port activation vs. calling a function or returning a value
 - PDU vs. actual parameter value

Network node model

- Each node has
 - A set of port types
 - A set of ports
- Each port type has
 - Message types
- Each message type has
 - Message format
- Each port has
 - A message queue (may be of length 1)
 - The queue may be shared

Input for architects and designers

- Requirements from:
 - Textual protocol specification
 - Customer stories – explained by domain experts

Elements of a protocol specification

- What to look for in the requirements?
 - Services of the system, module, component
 - Assumptions about the execution environment
 - Messages
 - Message formats
 - Behavior rules

Environment

- What to look for?
 - Is it reliable? → If not, then connection oriented protocol must be used and data must be protected with a checksum
 - Is it secure? → If not, then security and policy extensions may be necessary
- Differences between network node and software component
 - In software the environment is reliable, in communications it is nondeterministic
 - In communications the frequency of environmental events is higher

Service

- Input for both architects and designers
- What to look for?
 - Node types → This defines the network architecture
 - Is it reliable? → If yes, then flow control must be adopted
 - Is it synchronous? → If not, then message identification is necessary
 - Is it symmetric? → If yes, then P2P, if not, then client-server or SOA
 - Is it between neighbors? → If not, then addressing is necessary

Service in networking environment

- Specific to telecommunications architecture
- Service is subdivided into:
 - Data plane (Forwarding plane)
 - Control plane
 - Management plane
- Often different protocols implement the planes

Service in networking environment

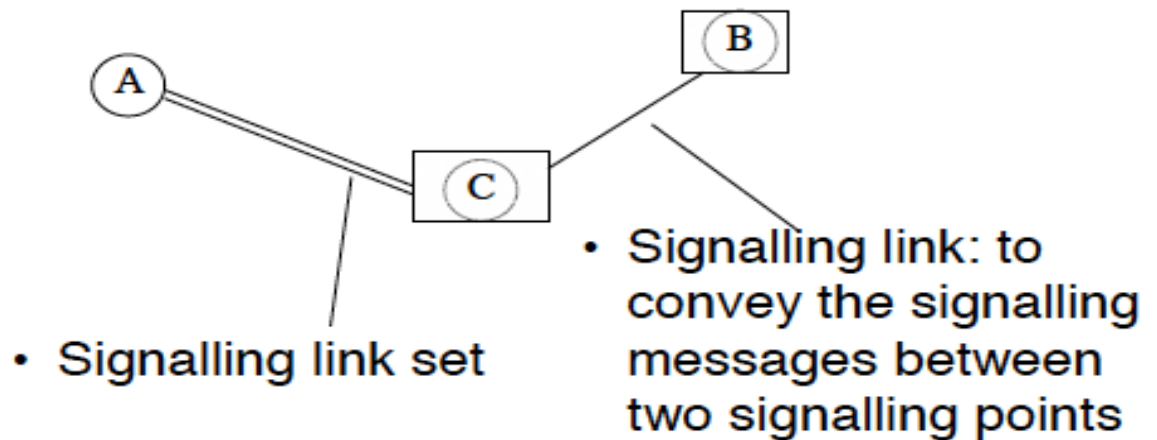
- Data plane
 - Forwards incoming requests, user traffic
 - Parses protocol headers
 - Queueing, QoS, filtering
- Control plane
 - Learns what the node should do with incoming requests -- control traffic
 - Maintains resource state and exchanges status information between peers
- Management plane
 - Controls and observes node status

Service examples

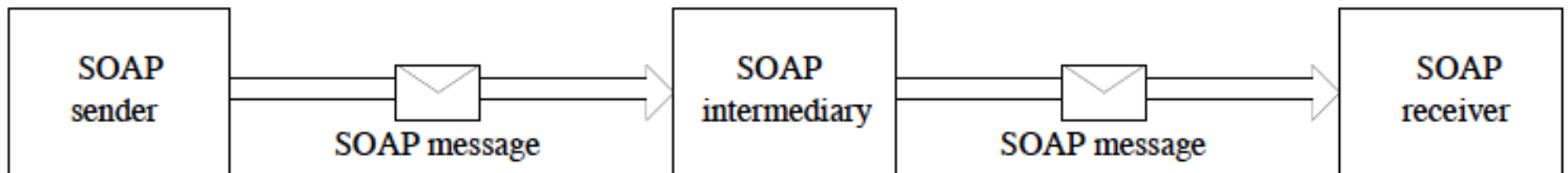
- SS7 MTP2
 - Historic control plane protocol for ISDN networks
 - Service: reliable transfer and delivery of signaling information between neighboring signaling nodes
- SOAP
 - Core protocol of XML web services (data plane, application layer)
 - Service: transfer of XML documents between neighboring SOAP nodes over a network

Network architecture example

- Network nodes may be on the same computer, even in the same software component
- SS7 MTP2



- SOAP

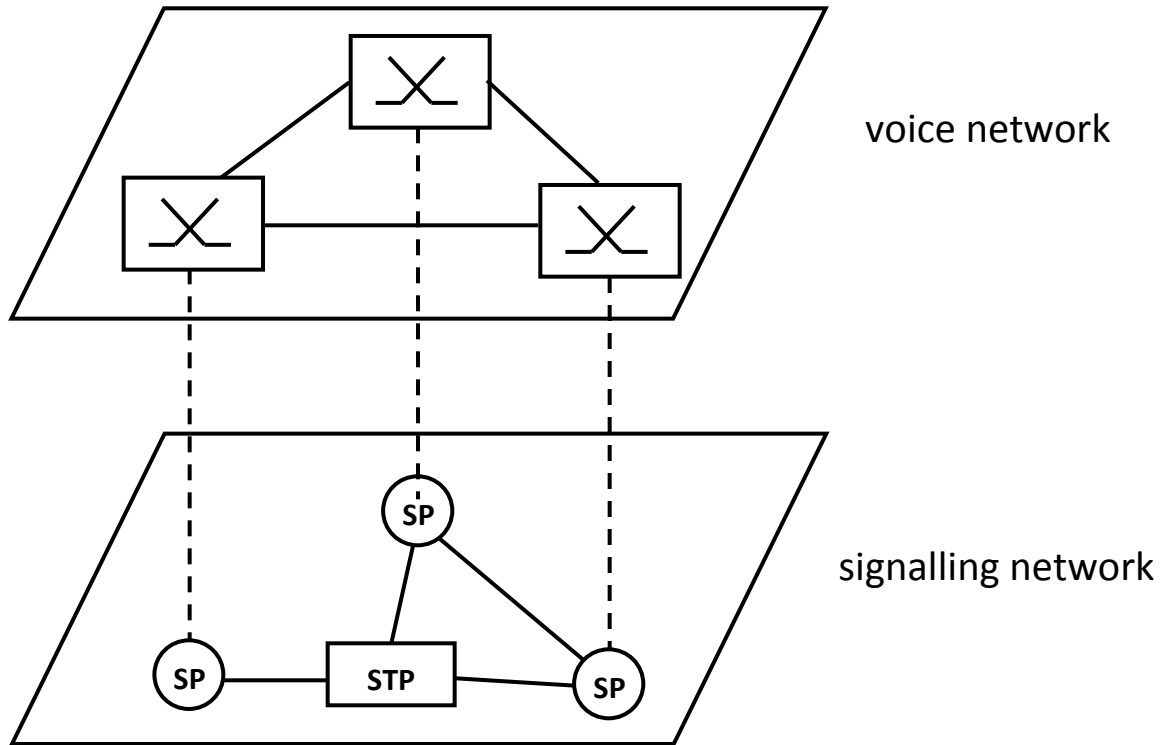


Multilayer and multitier

- Multilayer architecture
 - In communication networks
 - Usually between symmetric peers
 - Service is functionally decomposed into layers
 - Each layer has a small, well-defined task
 - Each layer has a unique language, a small set of functions
 - Higher layers are transparent to lower layers
 - Each layer has an overlay network
 - Layer protocols are components and hence replaceable

Multilayer and multitier

- Overlay network

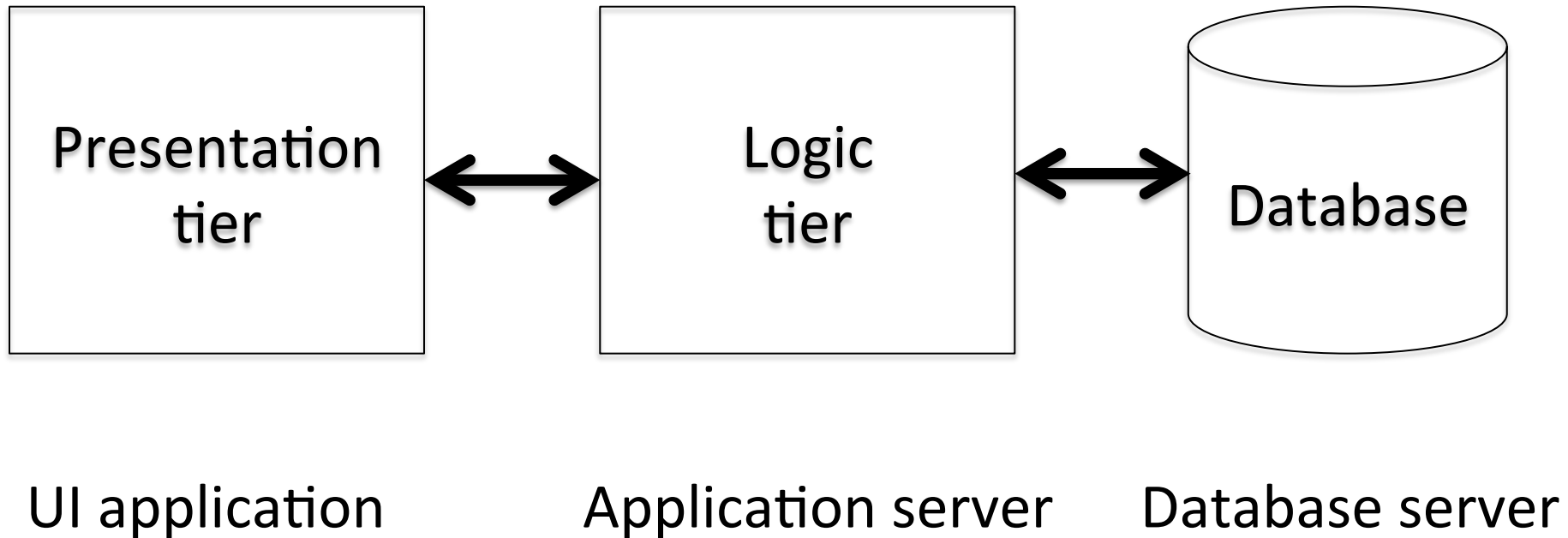


Multilayer and multitier

- Multitier (or N-tier) architecture
 - In software systems
 - Usually in client-server communication
 - Usually each tier is a separate network/
infrastructure node
 - Internal data representation
 - Unbounded number of data access functions
 - Data is not encapsulated but translated on tier borders

Multilayer and multitier

- Three-tier architecture

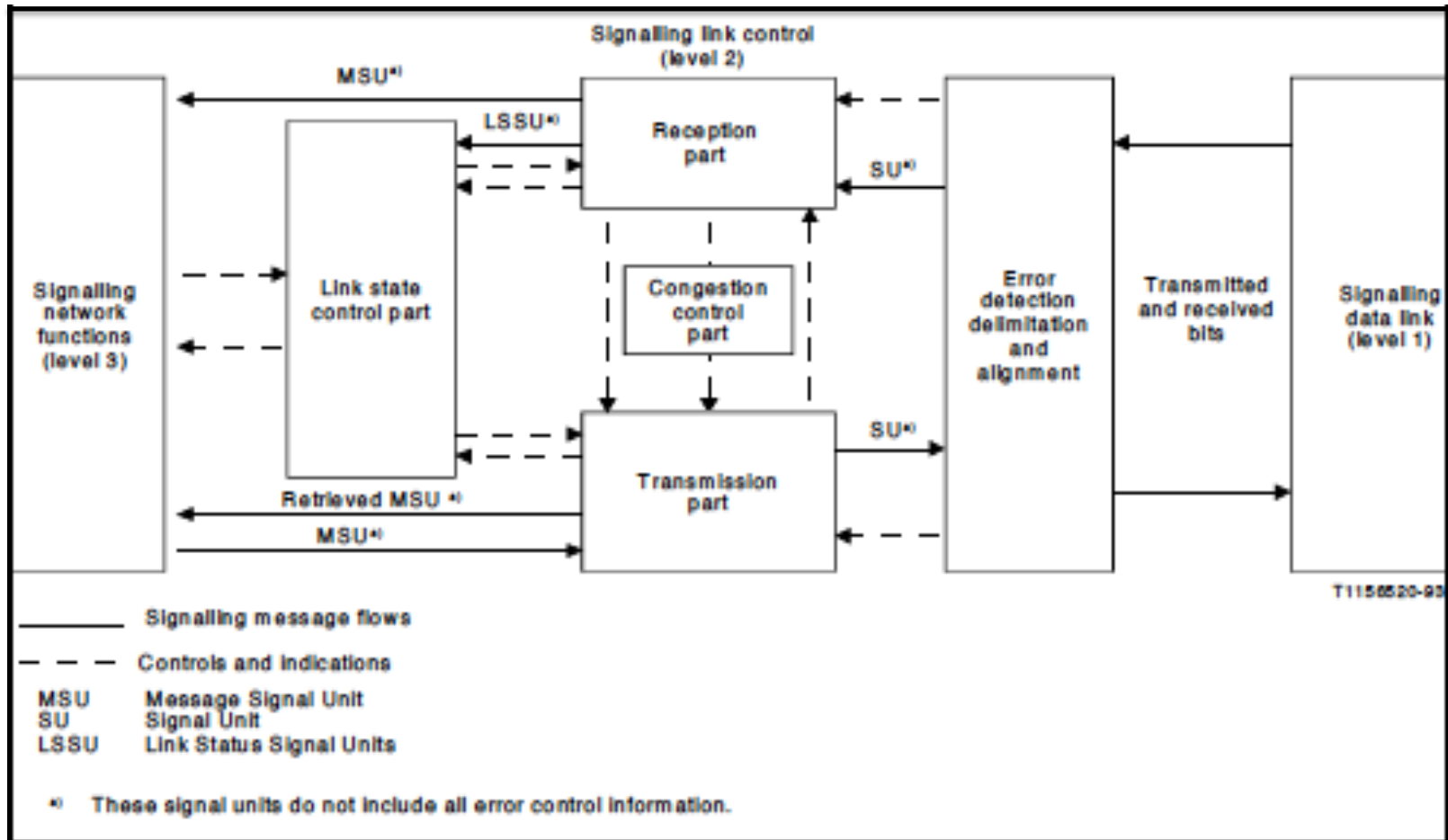


Multilayer and multitier

- Layer/tier elements
 - Client/transmitter
 - Server/receiver
 - Encoder-decoder → data serialization
 - Control, management
- If P2P, same transmitter and receiver in all nodes
- If client-server or SOA, unique transmitter and receiver in each node

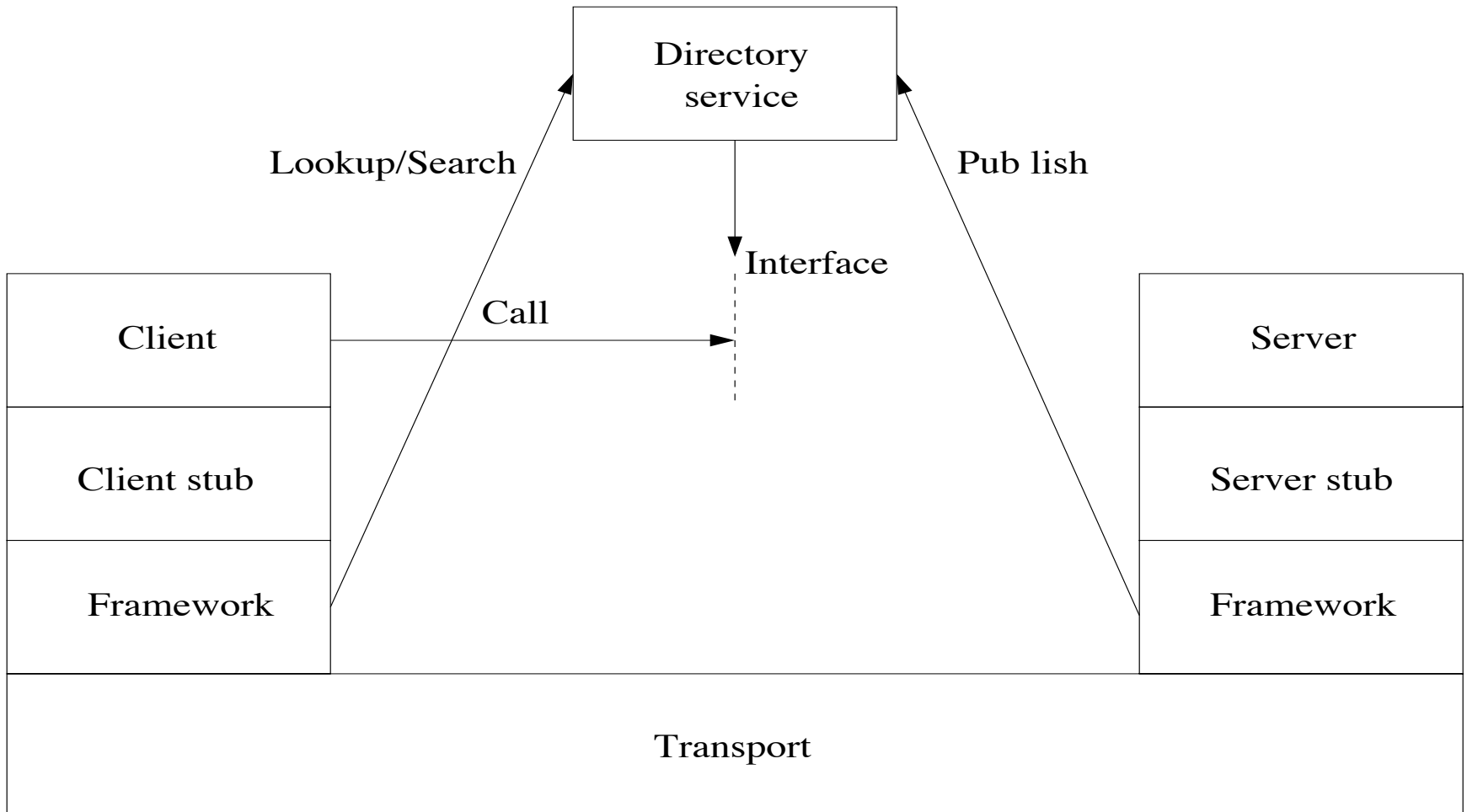
Example

- Internal structure of a telecom protocol



Example

- The CORBA distributed communications model



Message type

- Task of the architect
- Message = function identifier on an interface
- What to decide?
 - Synchronous or asynchronous
 - One-way, request-response, solicitation-acknowledgement, notification
 - What errors can occur?
- Functional decomposition: subservices → layered model

Message type

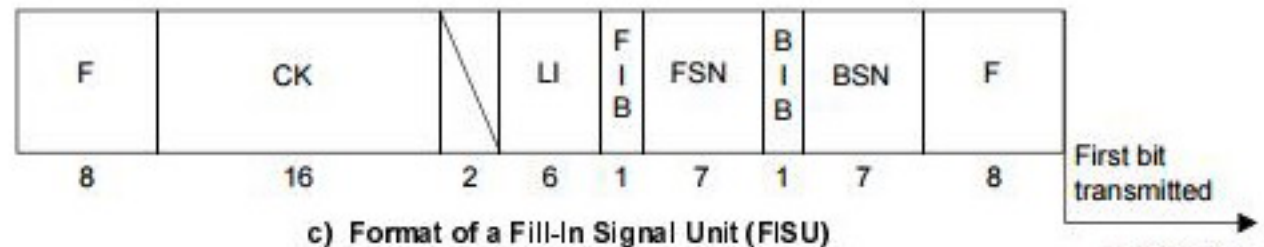
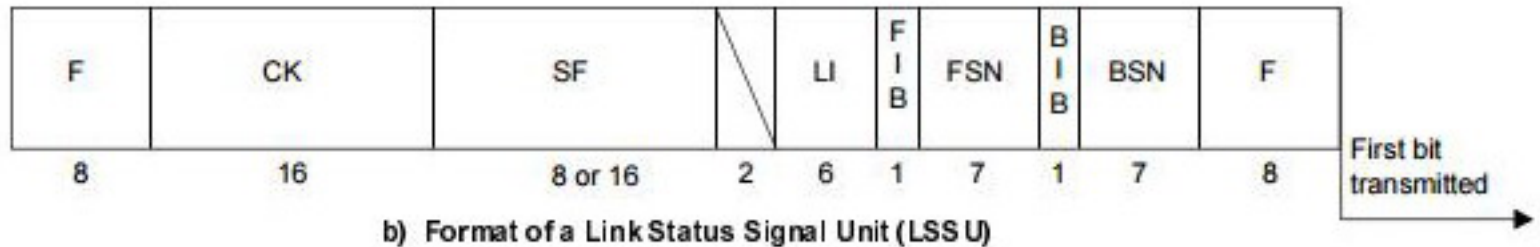
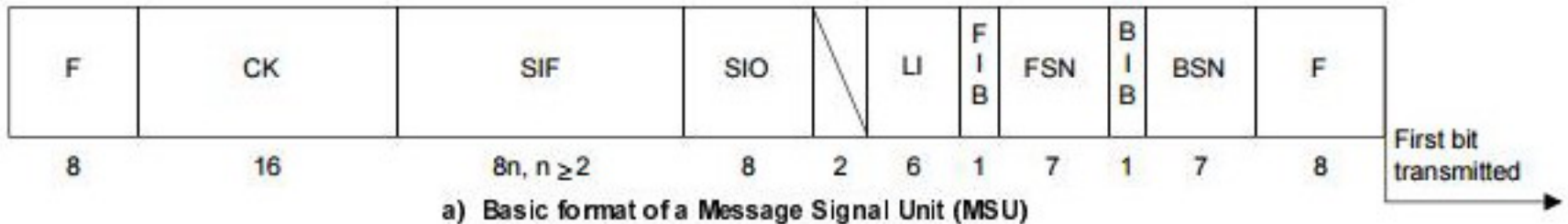
- How many message types should be defined?
 - As less as possible
 - All message types that appear on a public interface must be supported throughout the life cycle of the product – that can be expensive
 - Alternative: optional headers as key-value pairs

Message format

- Task of the architect
- Message format = function signature or return type
- What fields to include?
 - User data
 - Header data
 - Message number if asynchronous
 - Sequence numbers and fault and acknowledgement request indicators if reliable
 - Address if network of similar nodes exists

Example

- MTP2 message types and message format



Example

- SOAP message type and format

```
<?xml version="1.0"?>
<soap:Envelope
xmlns:soap="http://www.w3.org/2001/12/soap-envelope"
soap:encodingStyle="http://www.w3.org/2001/12/soap-encoding">
  <soap:Header>
    ... control data
  </soap:Header>
  <soap:Body>
    ... user data
    <soap:Fault>
      </soap:Fault>
    </soap:Body>
</soap:Envelope>
```

Message queue

- Each port of a component is associated with a queue
- Queue types:
 - Single FIFO – ordering of arrivals are lost
 - Multiple (per port) FIFO – which input queue to consider next
 - Typed message pool – which type of message to look for in the queue
 - General message pool – reception may depend on message parameters (TTCN-3 – later)
- A delaying network link itself is a queue

Behavior modeling

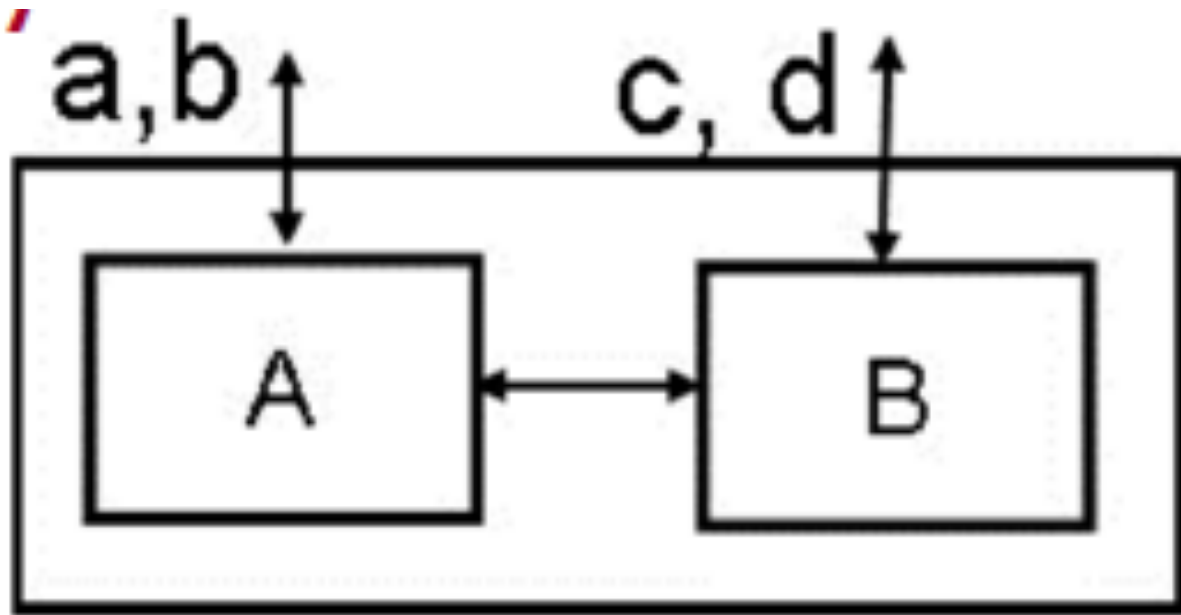
- Mealy finite state machine
 - Set of input events
 - Incoming port activations
 - Internal timeouts – not observable from the outside
 - Set of output events
 - Outgoing port activations
 - Set of states
 - Variables dedicated to remember I/O event history
 - Not observable from the outside
 - State transition rules
 - What is the next state and the output event upon an input event in a certain state

Problems from the environment

- Transport medium is not reliable
 - modeled as a non-deterministic function
- Messages may be
 - delayed
 - lost, corrupted
 - inserted, duplicated
 - reordered

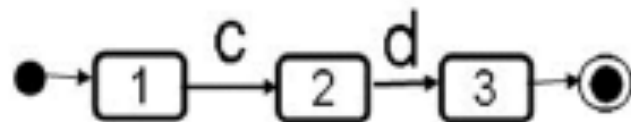
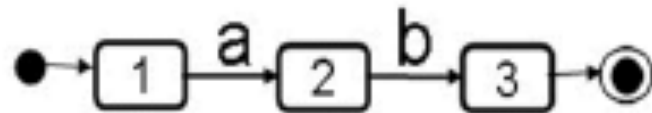
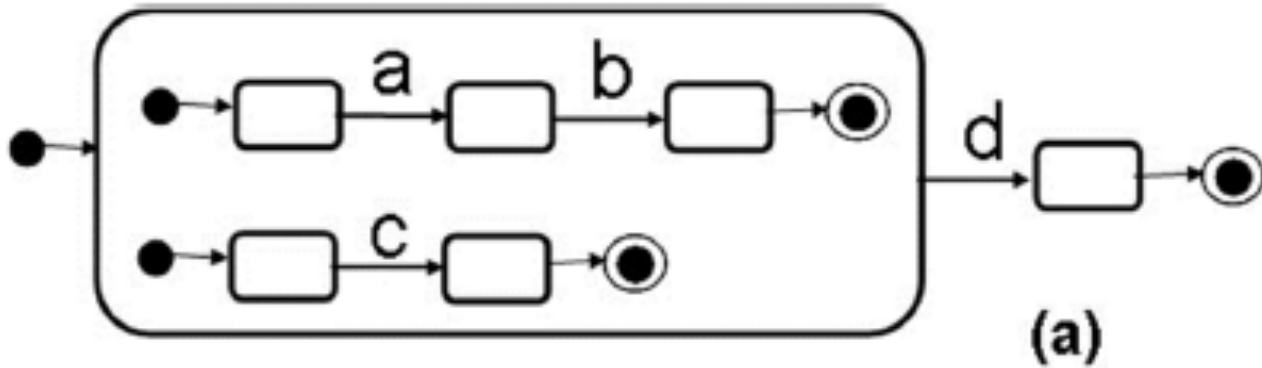
Example

- Protocol layer with two ports
- Components A and B may be local or distributed



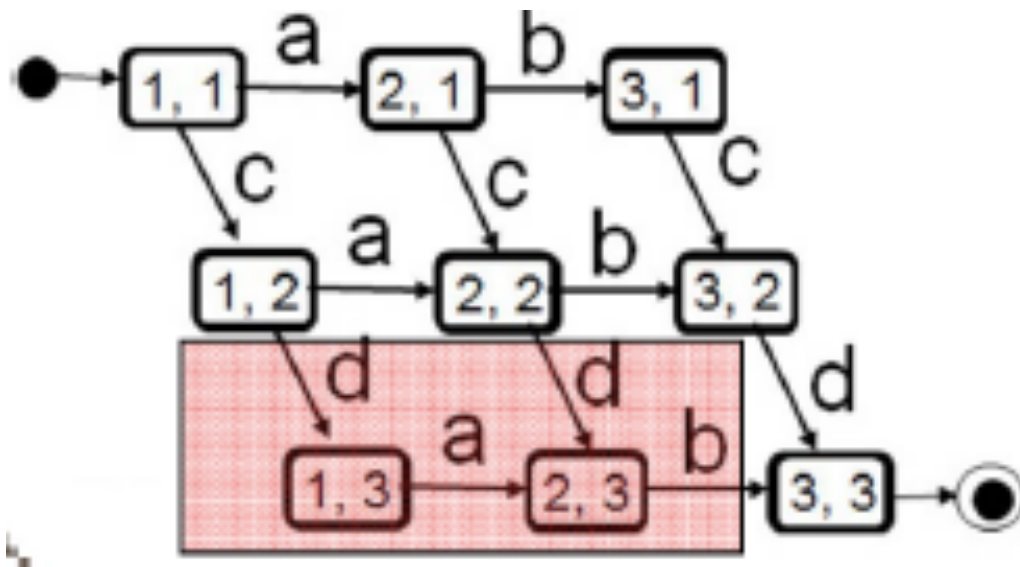
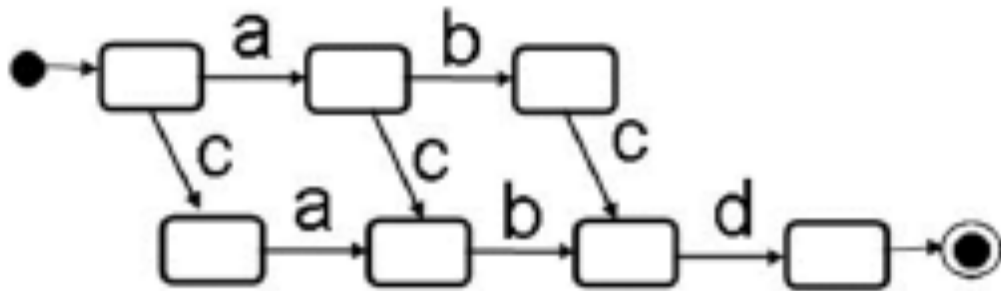
Example

- A behavior expressed with state machine



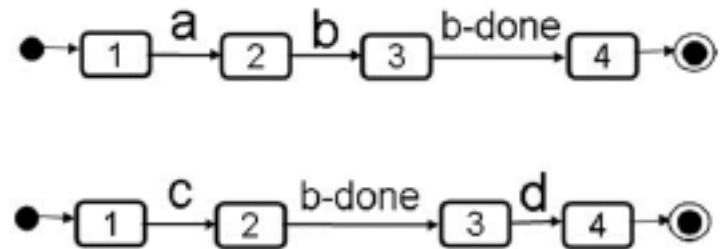
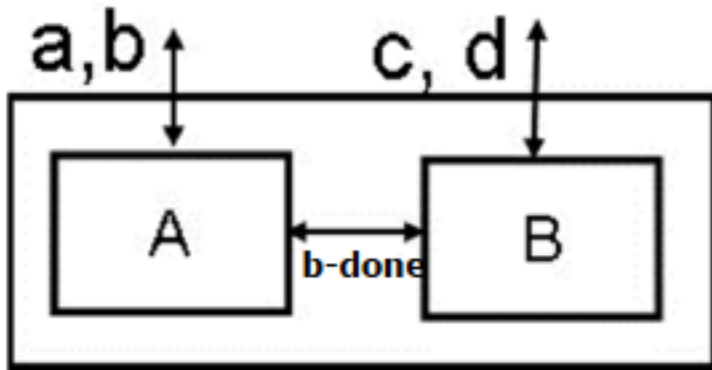
Example

- Desired and invalid behaviors



Example

- How to fix?
 - Rendezvous communication
 - Avoids cross-over of messages over the interface
- If the synchronization message is lost, the system deadlocks, hence must be protected with timers



Synchronous, asynchronous

- Rendezvous communication
 - Synchronous
 - Remote procedure call
 - Caller must wait for the called – return
 - Queue of one element
- Message passing communication
 - Non-blocking
 - Multiple messages can be on route – queuing
 - Callback instead of return type
 - Message identification necessary

Addressing

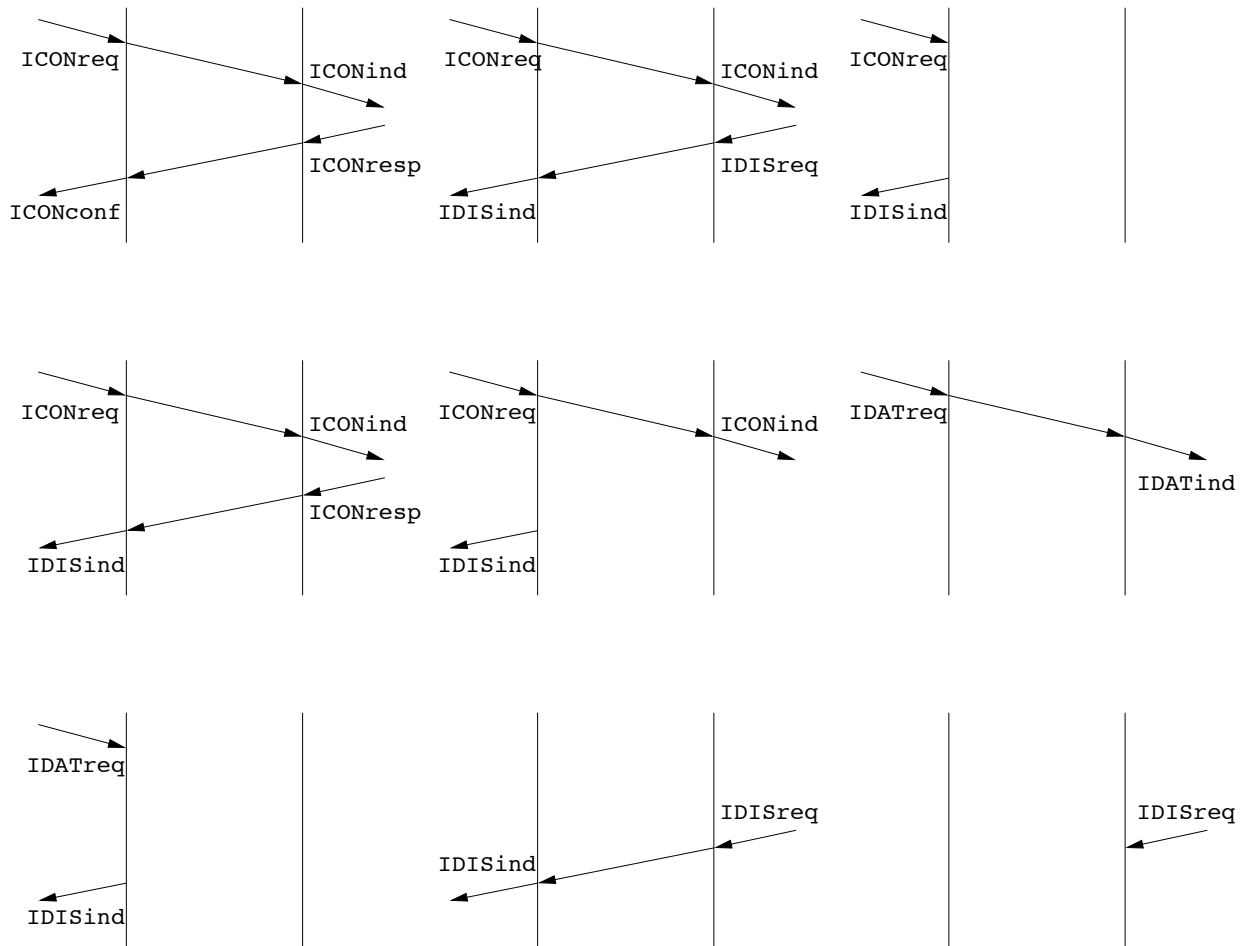
- Endpoint addressing is necessary if the communication is not P2P
 - Communication is asynchronous
 - The request P2P session is terminated once delivered, and the peer must know where to send the callback message
 - Next-hop routing is used
 - Over a link layer service or WS-Addressing
 - Each node along the path examines the destination endpoint address, and forwards the message to a node closer to the destination

Connectionless, connection oriented

- If messages must arrive exactly or at least once or in order a P2P connection must be set up between the two nodes
- Connectionless
 - Example: HTTP
- Connection oriented
 - Connection/session id
 - Tracking delivery of messages
 - Extra message types for connection setup, termination and management

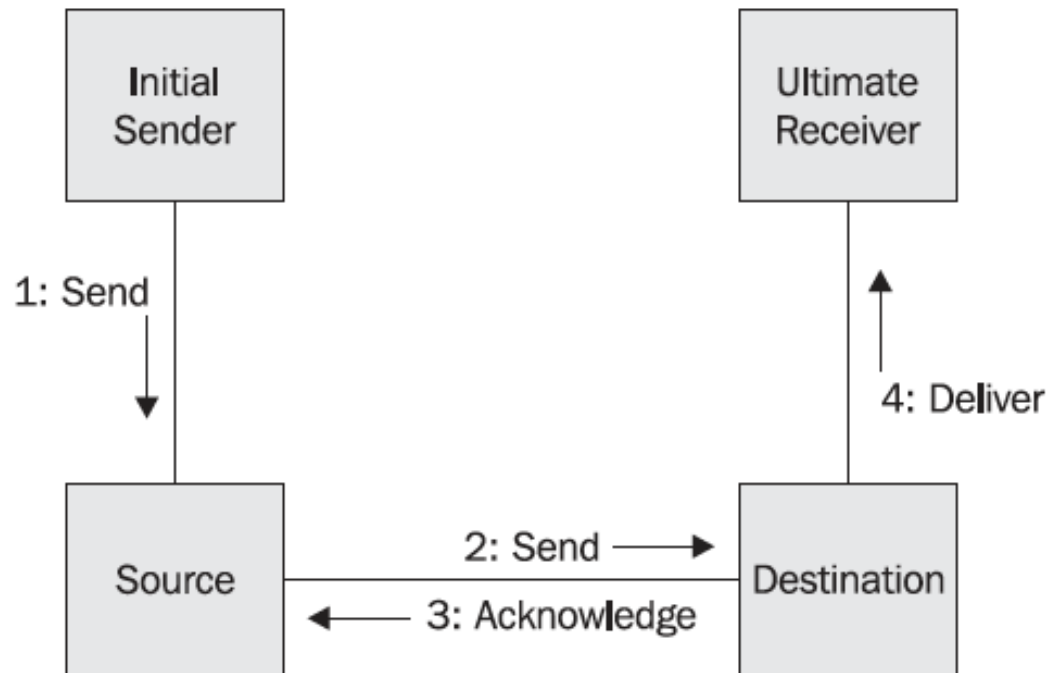
Example

- Connection orientation with primitives



Example

- SOAP is unreliable: WS-ReliableMessaging extension

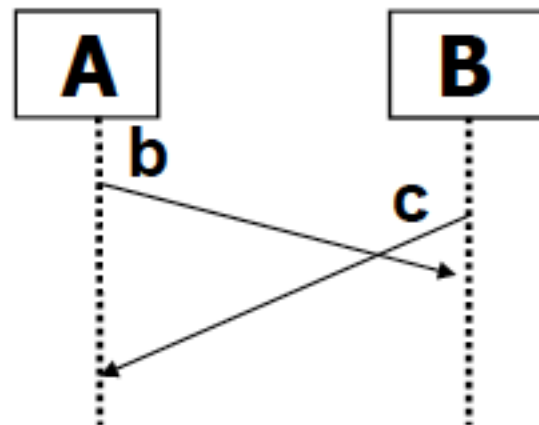
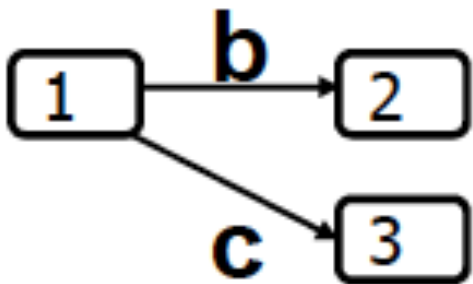


Example

- SOAP connection control
 1. CreateSequence ($s \rightarrow d$)
 2. CreateSequenceResponse(id) ($d \rightarrow s$)
 3. Sequence(id, msgId) ($s \rightarrow d$)
 4. SequenceAcknowledgement – optional ($d \rightarrow s$)
 5. AckRequested ($s \rightarrow d$)
 6. TerminateSequence(id) ($s \rightarrow d$)
- Connections and messages have id

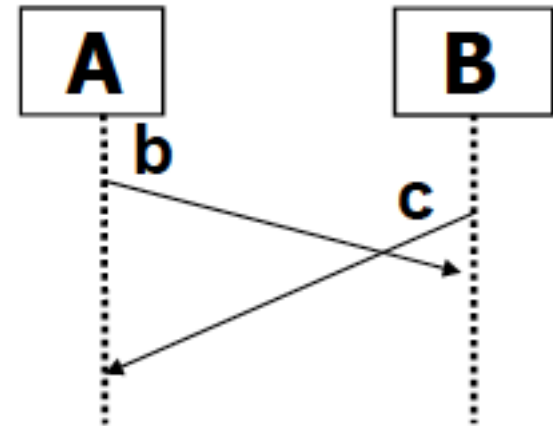
Racing messages

- If the source of *b* and *c* is the same component, the ordering is a local choice
- If different component initiate *b* and *c*, then those are competing



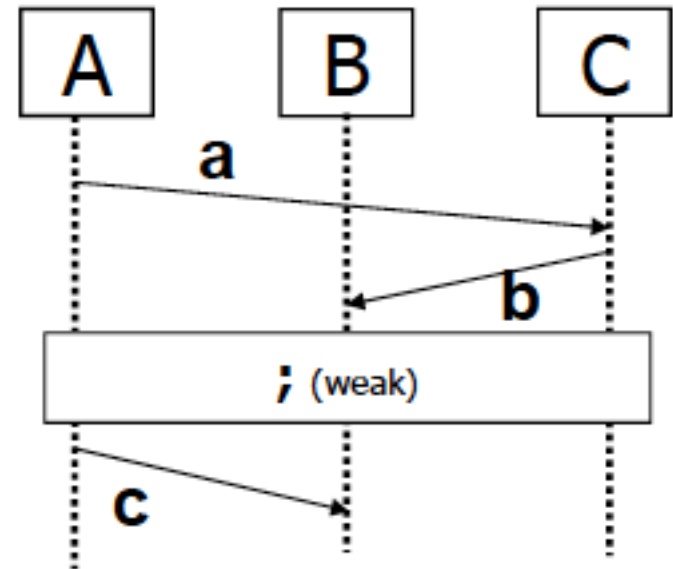
Racing messages

- Both components will be in different states
- Possible solutions:
 - Coordination protocol
 - Priority for one side
 - Non-specified reception
 - Undefined behavior
 - Error state
 - Message is ignored



Racing messages

- Two messages sent from a component arrive at different ports of another component
- Ordering enforced in A
- That has no influence on B
- Coordination is necessary
- Problem a b is a loop
 - When does the loop end?

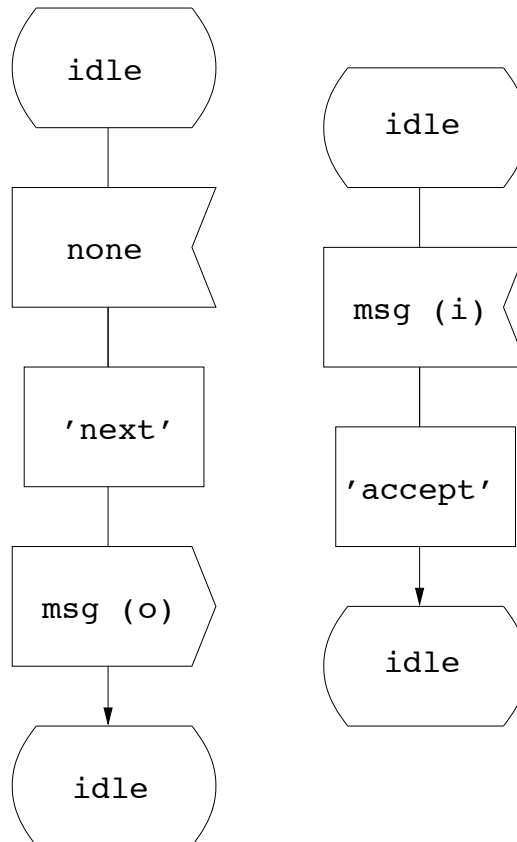


Flow control

- The sender produces data faster than the receiver could process them
- Protects against: message deletion, insertion, duplication and reordering, racing conditions, deadlock
- Elements:
 - Sequence numbers
 - Control messages
 - Timers

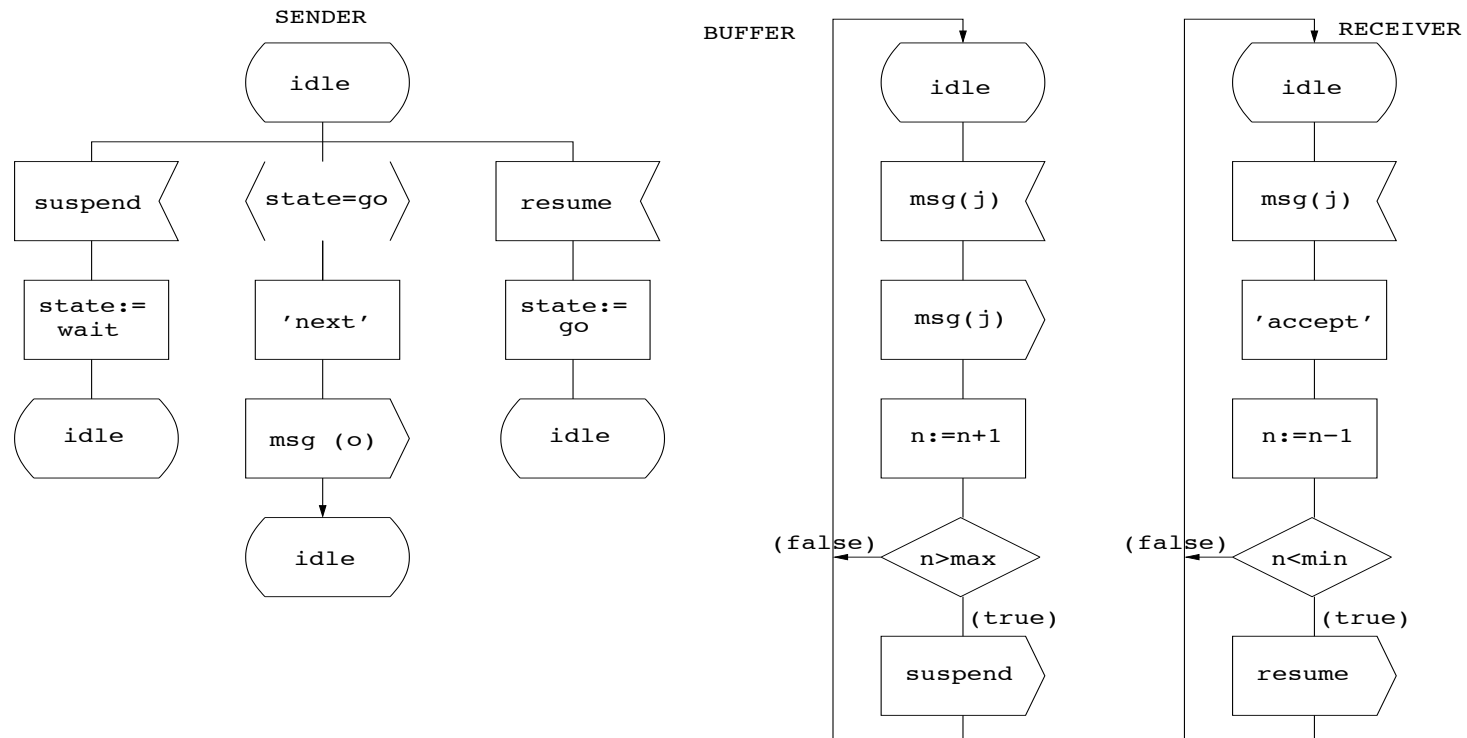
No flow control

- The receiver process must be faster than the sender



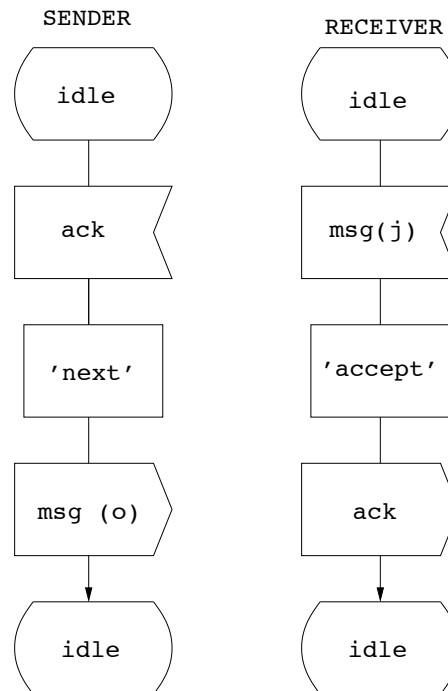
On/off

- The receiver sends suspend if its buffer is full and resume if it is empty
- Problem: loss or delay of control message



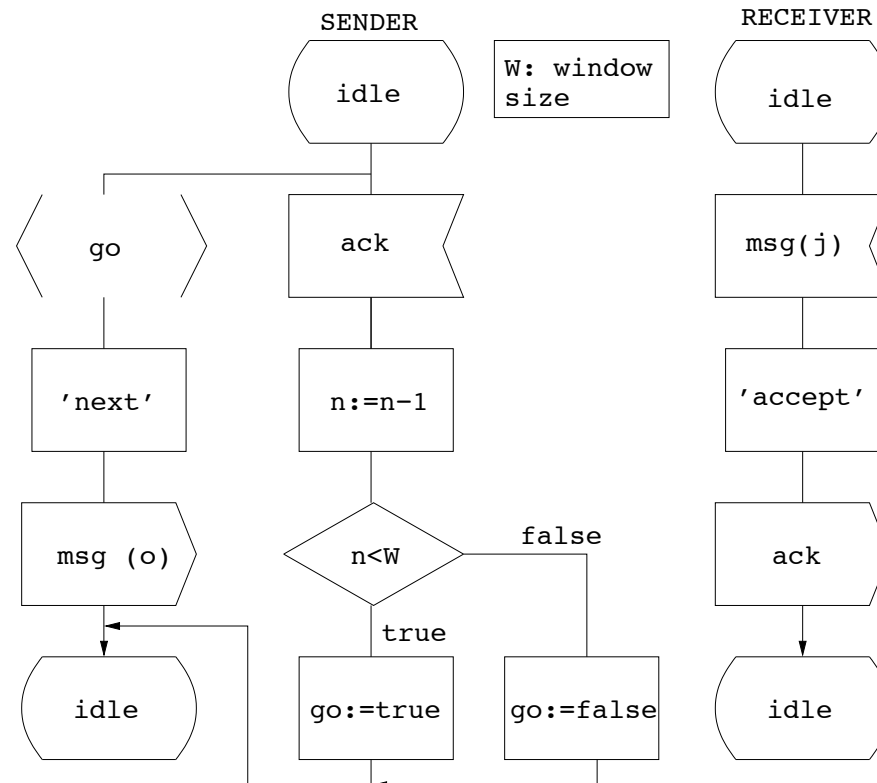
Stop and wait

- The receiver acknowledges each message, the sender must wait for the acknowledgement
- Deadlock if a control message is lost



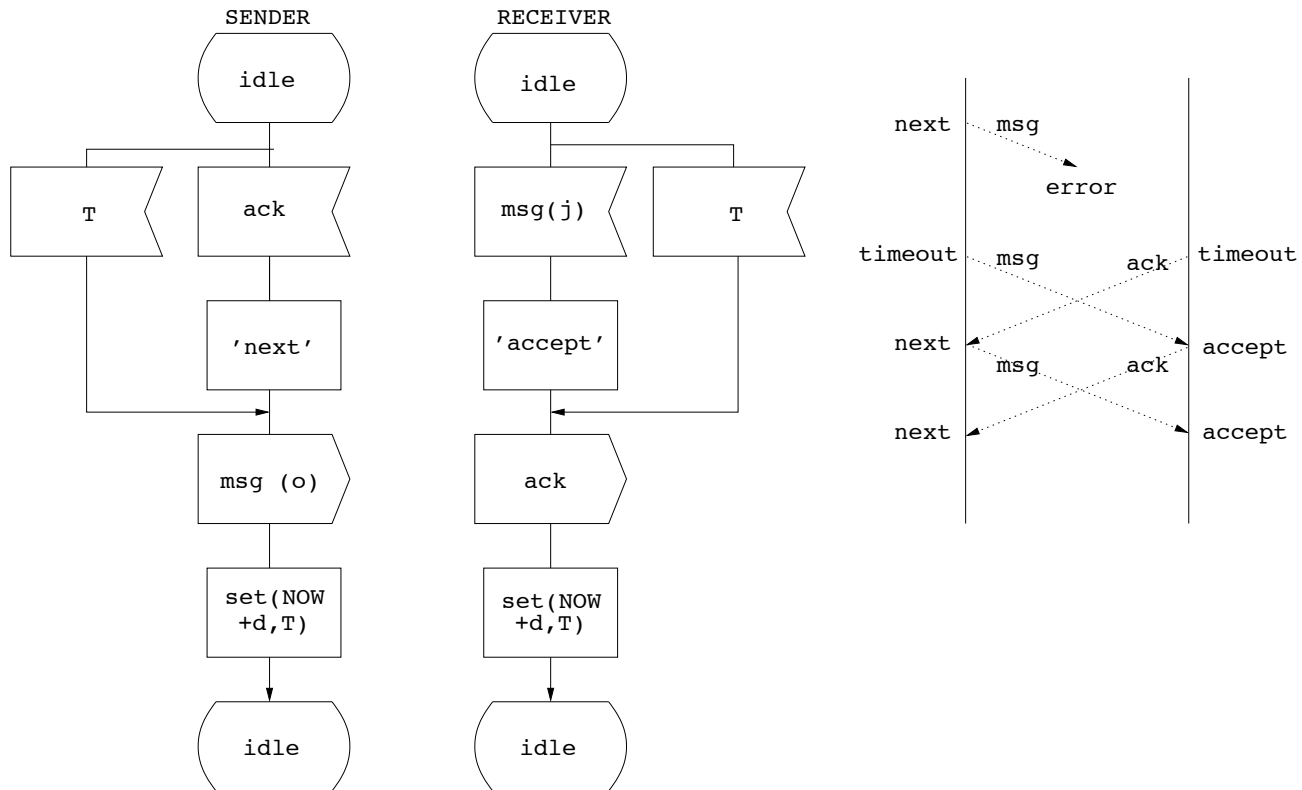
Windowing

- The receiver acknowledges each message, the sender knows the buffer size of the receiver
- Problem: lost, inserted reordered, duplicated messages



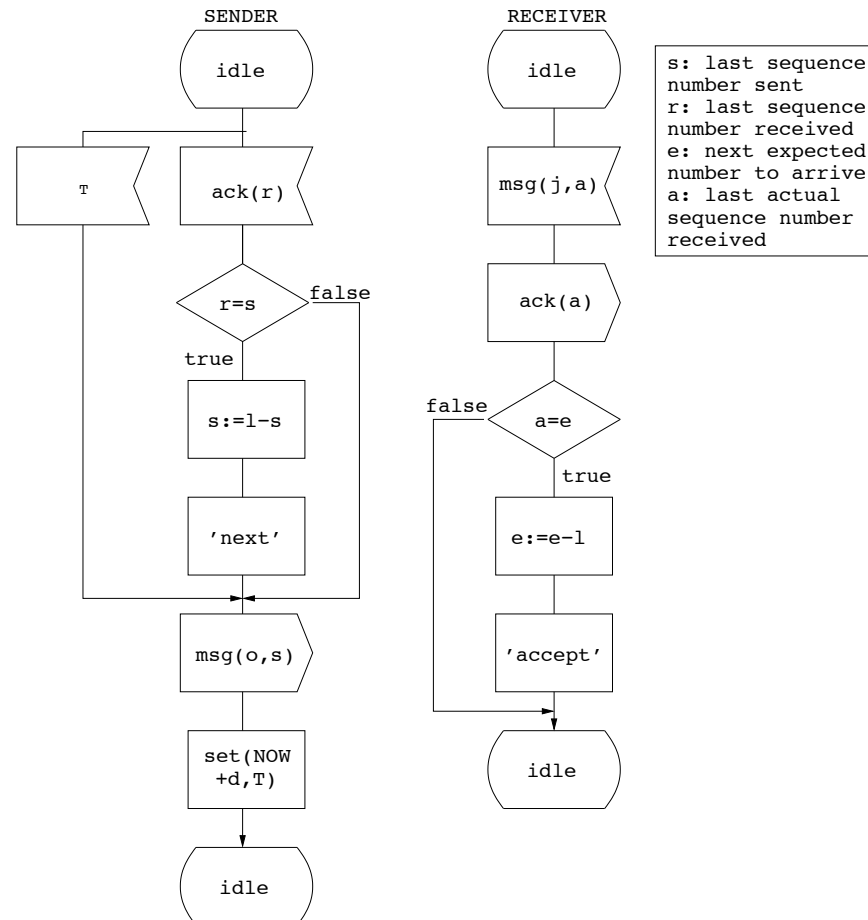
Timeout

- Protection against message loss
- Only one of the processes may initiate the retransmission



Sequence numbering

- Provides retransmission of lost messages



Sliding window

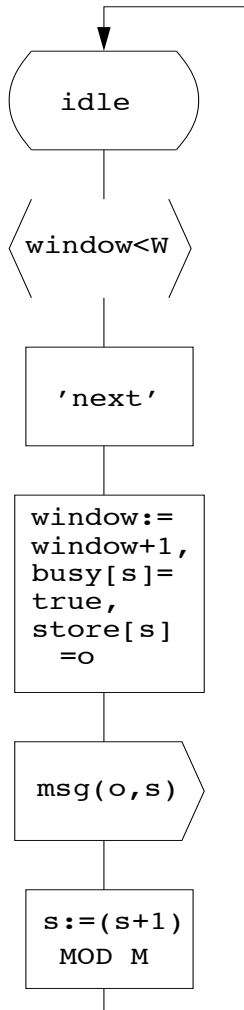
- Provides the retransmission of lost messages
- Elements:
 - Larger sequence numbers, total M
 - Window size: $W \leq M/2$
 - Sent but not yet ackd messages are stored on the sender side
 - Received but not yet accepted messages are stored on the receiver side
 - Receiver side buffer

Sliding window

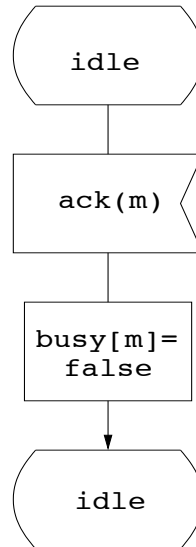
- The window:
 - Next message to be sent (s)
 - Number of unacknowledged messages (window)
 - Oldest unacknowledged message (n)
 - Latest unacknowledged message (m)
 - Next message to accept (p)

Sliding window

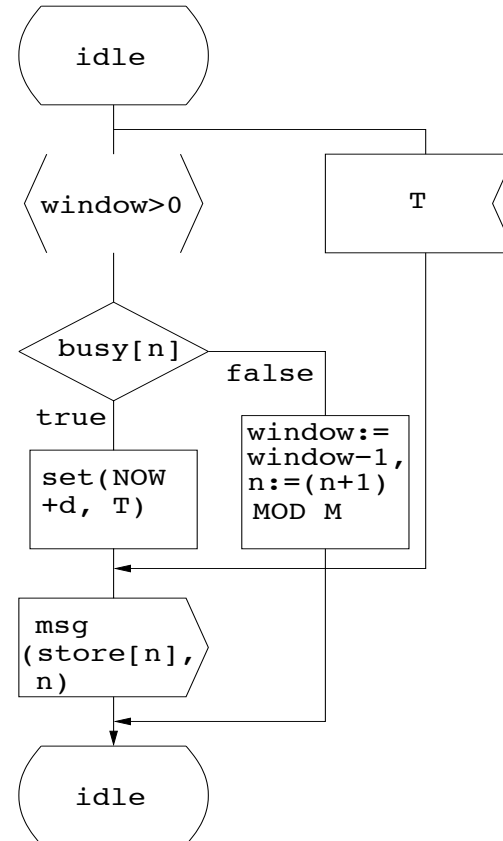
Transmission Process



Acknowledgement Process

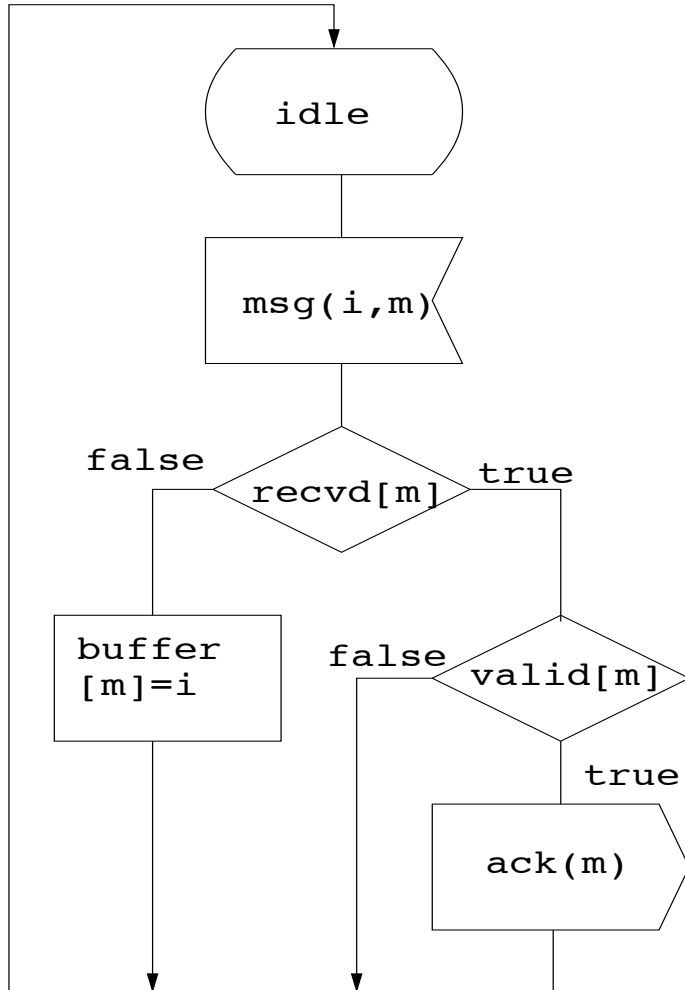


Retransmission Process

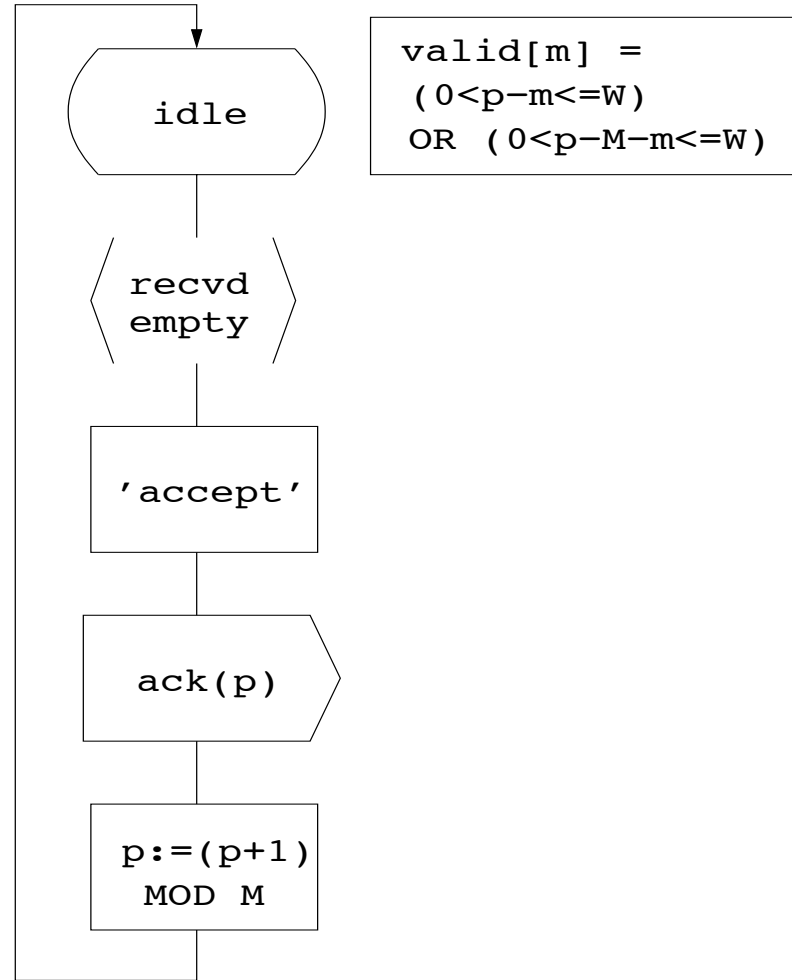


Sliding window

Receiver Process



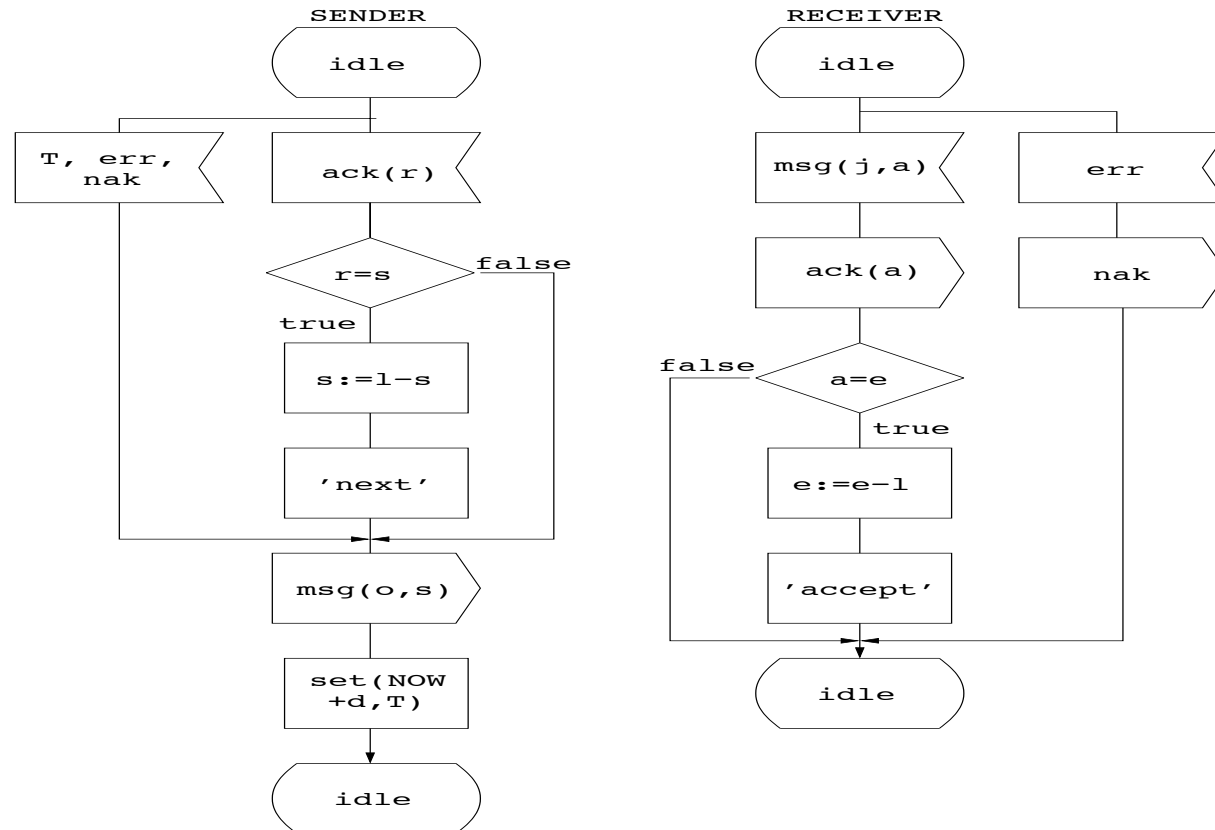
Accept Process



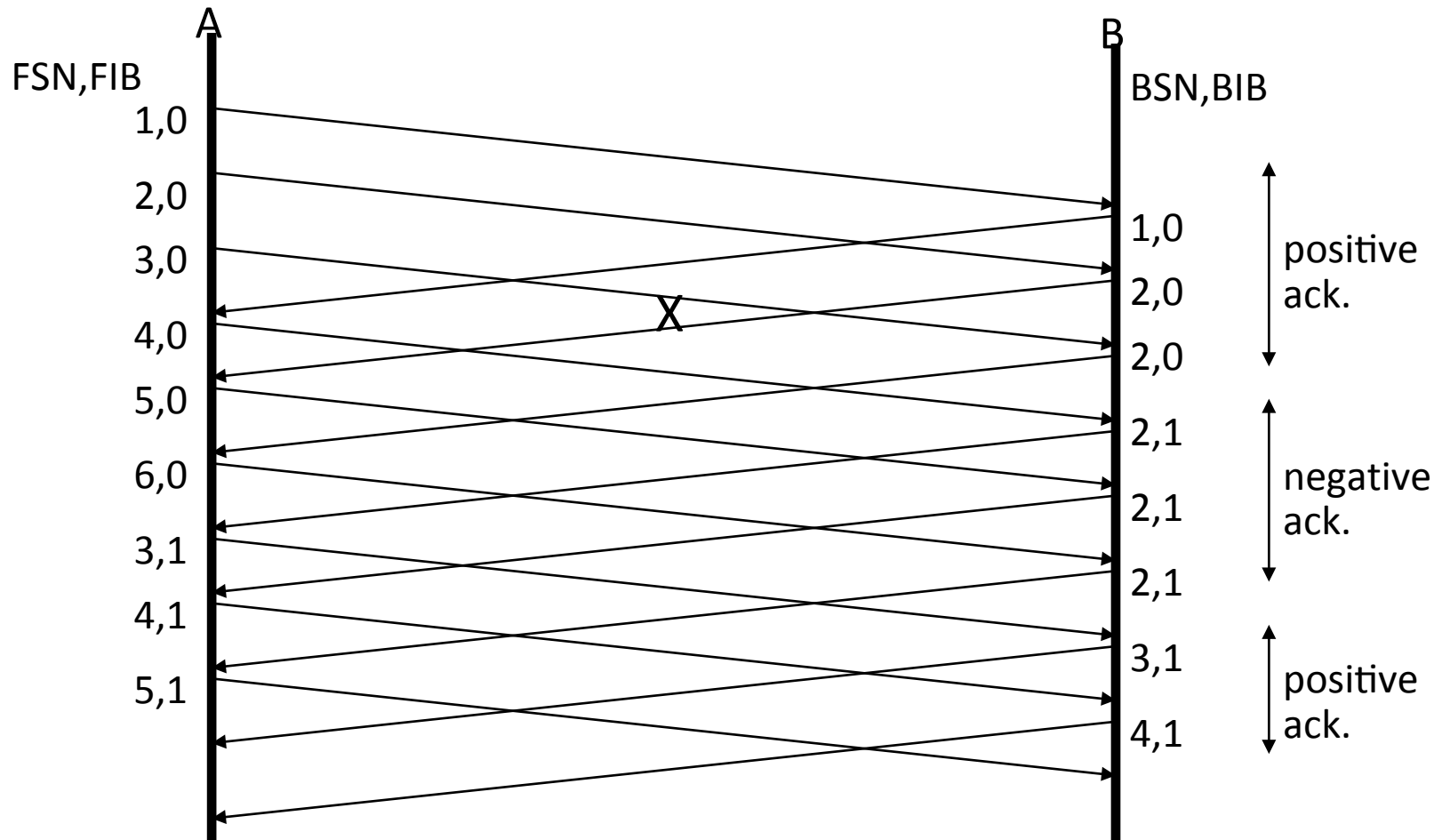
`valid[m] =`
`(0 < p - m <= W)`
OR `(0 < p - M - m <= W)`

Negative acknowledgement

- The receiver forces the retransmission of damaged messages instead of sender timeouts
- Block acknowledgement



Example



Policy

- Rules for participating in communication sessions
- Must be agreed upon before the communication takes place
- Policy consists of:
 - Subject
 - Domain
 - Rules, assertions
 - Required, optional, observed etc.
 - Decision points
- Examples:
 - In networking: SNMP
 - In software: Java Security Policy or WS-Policy