

Sensor networks and applications

Network topology, routing

Topology

 Tipically, sensor networks are self-organizing, ad-hoc networks, the network structure "can not be planned".

 The physical connections (links) are randomly created, resulting in a random (physical) topology.

- The <u>logical topology</u> construction can be important!
 - Taking into account <u>scalability</u>.



Distributed (flat) structure

- There is no constructed logical structure, all nodes take part in network control.
- Since the nodes have information only about their neighbors, this information must be distributed within the whole network from time to time.
 - E.g., every node broadcasts its routing table periodically.
- Drawback: not scalable! (e.g., for thousands of nodes...)



Clustering

- The network is divided into so-called clusters, so that each node belong to at least one cluster.
- All clusters have a cluster head, which controls all the nodes within its cluster.
- Gateway nodes proved communication between neighboring clusters.
- Scales better.
- Problem: All clusters know the neighboring clusters, but how can it be informed by the far away clusters?
 - The same problem on a higher hierarchical level!
- Solution: Cluster heads are organized into a hierarchic tree.



Clustering (cont'd)

- Optimal clustering is not easy
 - "NP-hard", but there are $O(n^2)$ heuristics

- Another problem is that theload on the cluster heads can be much higher than on ordinary nodes.
 - Solution: The cluster head role is changed from time to time.



Next: Routing

- Routing planning issues
- Modeling the network
- Structure-based solutions
 - flat
 - hierarchic
 - location-based



Content

- Medium Access Control
 - S-MAC

- Networking layer
 - WSN topology
 - Routing

- High number of network nodes:
 - Global addressing is not possible, because unique ID can not be provided to all the nodes.
 - Consequence: The "traditional" IP-based protocols can not be applied.

- Sensors are deployed in an ad-hoc way:
 - The network must be self-organizing.
 - Random node distribution must be dealt with.
 - Self-managing is required.



- Typically, sensor nodes are static after deployment.
 - In contrast to mobile ad-hoc networks, where the nodes can move freely.
 - Some nodes can be mobile according to the given application (but with low mobility, typically).

- Typically, the information is flowing from many sources (sensors) to a single node (base station).
 - It can be multicast or peer-to-peer as well!



- Sensors have limited energy, computing- and storage capacity.
 - Efficient recource-management is needed.

- Sensor networks are application specific.
 - The design requirements are different for different application areas.
 (E.g., high precision monitoring systems vs. periodic wheather monitoring.)

- The location aware operation is important, because data collection is location dependent.
 - Localization based on triangulation using reference points.
- The collected data are reduntand.
 - E.g., more sensors measure and collect data about the same phenomenon.
 - Redundance can be exploited during routing.
- Most WSNs are data centric.
 - We are looking for data that are attribute-based. (E.g., where is the temperature higher than 40 degrees?)



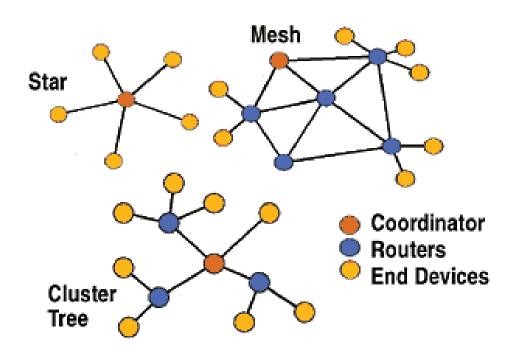
Next: Routing

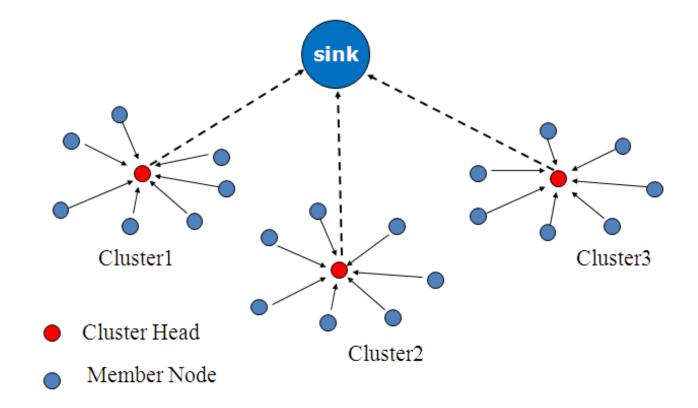
- Routing design issues
- Network and routing modeling
- Network topology-based solutions
 - Flat
 - Hierarchic
 - Location based



Routing

WSN topology – clustering







- Csomópontok telepítése: deterministic or random
 - <u>Deterministic</u>: nodes are deployed manually, the data traffic is carried by routes that are set up in advance.
 - Random: the infrastructure is set up in an ad-hoc way.
- Energy efficiency without loosing reliability and accuracy.
 - Dual role of nodes: data source and <u>routing</u>.



Data model

Information sensing and data transmission is application specific.

Sensing and data transmission can be

- time-driven (continuous)
 - e.g., periodic monitoring
 - Sensrs become active from time to time, sense and transmit data.

event-driven

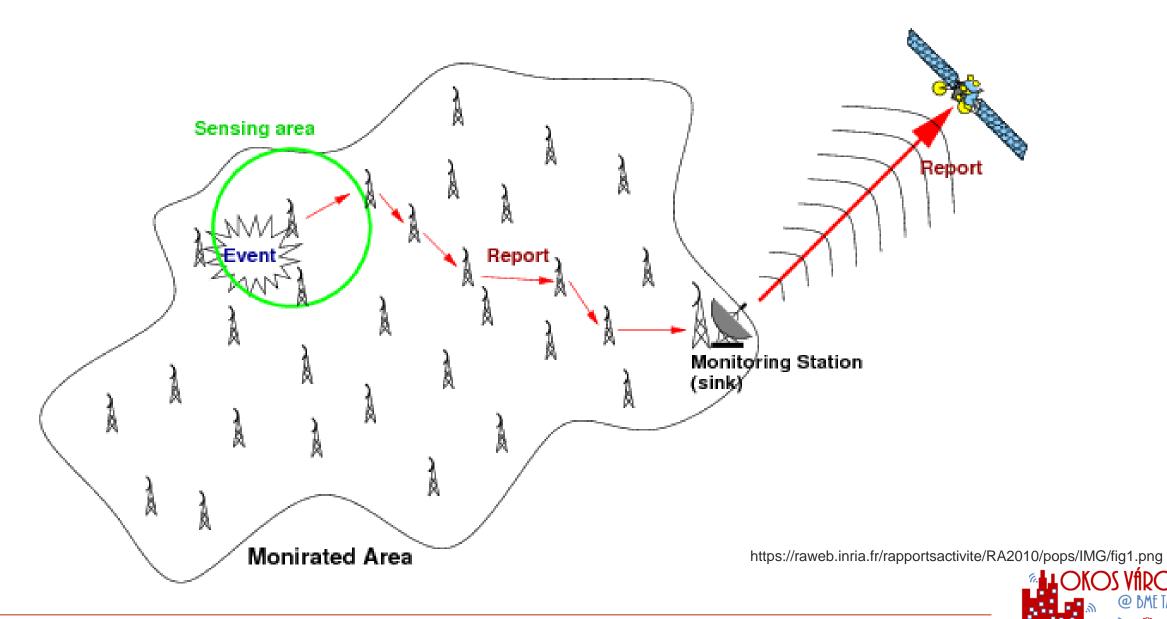
- Sensors react to sudden changes in the environment.
- can be important in time criticl applications.

query-driven

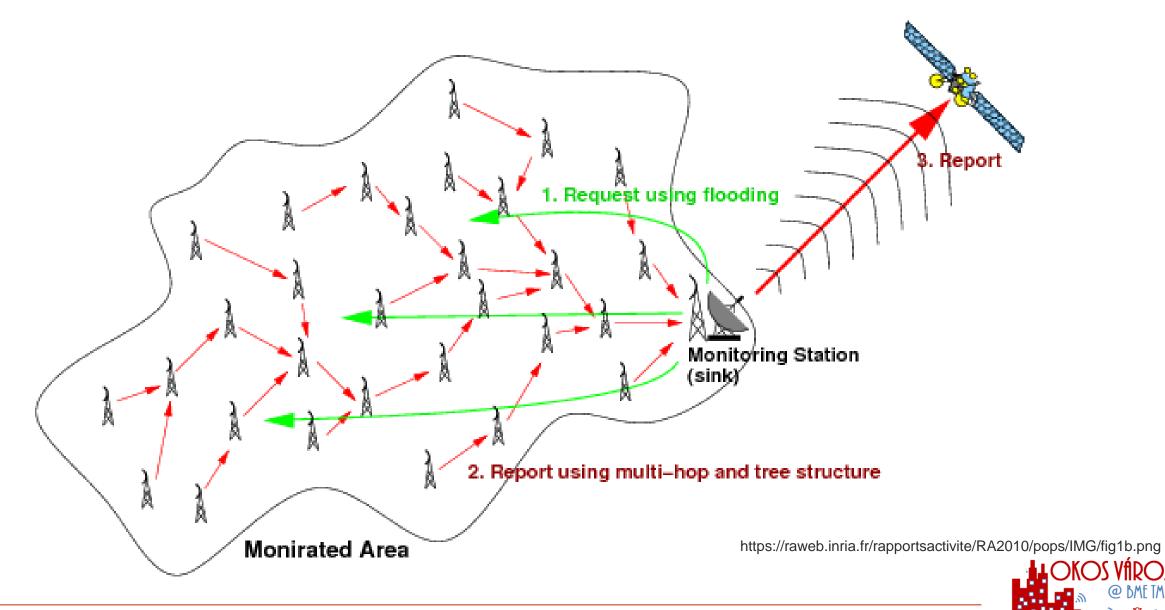
- Nodes answer to a query of a "controller"
- Routing heavily depends on the data model used!



event-driven network (example)



Query-driven network (example)



- Node/link heterogeneity
 - In most cases the sensor nodes are identical (homogeneous)
 - Afew nodes can be different according to the application.
 - E.g., they sense different physical things (video, noise, temperature).
 - The special sensors can be deployed differetly.
 - Data generation and transmission speeds can be different as well (e.g., temperature vs video).
 - Different data can require different QoS treatment.
 - E.g., cluster heads can be well positioned, stronger nodes with higher energy.



Robustness

- Any sensor node can drop out from the network for different reasons (e.g., interference, physical damage, energy depletion).
- Loosing a single node must not violate the operation of the whole application.
- The MAC and routing protocols must be able to deal with the loss of <u>more</u> nodes.
- Multi-level redundancy is necessary for building a sensor network.

Scalability

- The number of sensors in a WSN can be hundreds or thousands(!).
- E.g., as a consequence, the implemented routing table within a node can not contain all routing information to any other node.



Network dynamics

- In most of the cases the nodes are static, they don't move.
- For mobile nodes (e.g., mobile BS) the mobility handling can be a great challenge in routing.
 - E.g., the network topology changes in time, ...
- The sensed event itself can be static or dynamic as well. (e.g., temperature vs. target tracking)

Connectivity

- Typically, the degree of a network node is high due to the high node density.
- Connectivity can change in time (e.g., nodes deplete their energy sources).



Coverage

- The sensors can monitor only its close environment with adequate accuracy.
- The full area coverage (and keeping the coverage in time) is an important target in planning.

Data aggregation

- The sensors generate and send data with high redundancy.
 (E.g., temperature values at nearby locations)
- The amount of data to send can be significantly reduced by extracting the valuable information from raw data.
 - E.g., minimum, maximum, average
- By data aggregation higher energy efficiency can be reached.
- Signal processing techniques can alo be applied for data compression.



- Quality of Service (QoS)
 - E.g., the measured data should be sent within a time frame, otherwise it will be outdated.
 → For time critical applications the <u>delays must be bounded</u>.
 - For most of the applications, the energy efficiency is more immportant than small delay.
 - E.g., <u>Idea</u>: The data quality can be decreased linearly with energy depletion.
- Energy-aware routing is the key is WSNs!

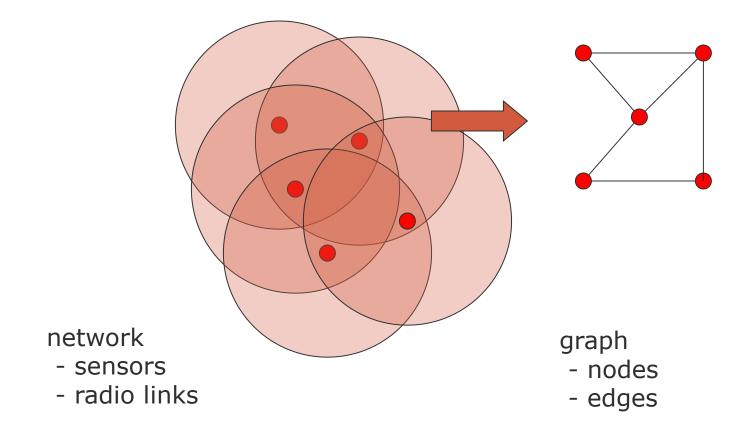


Network and routing modeling

Graphs, algorithms

Modeling network and routing

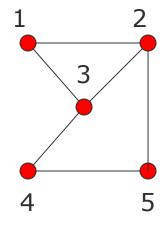
- Mathematical model:
 - network -> graph
 - routing algorithms -> graph theory algoritms





Modeling network and routing

- Network graph
 - Edges can be directed or undirected
 - E.g., a station does not take part in message relaying
 - The notion of distance can be used between any two nodes
 - E.g., the "cost" or "quality" of the link can be given
- Notations: G(N, A), where G is the graph and N is the set of nodes, A is the set of edges.



$$N=\{1,2,3,4,5\}$$

 $A=\{(1,2), (1,3), (2,3), (3,4), (2,4), (2,5)\}$

Network graph, definitions

path:

$$(n_1, n_2, ..., n_m), n_i \in N, (n_i, n_{i+1}) \in A$$

- route: Path where none of the edges are repeated.
- The graph is <u>connected</u>, if there exists a route from each node to any other node.
- <u>loop</u>: route, where $n_1 = n_m$
- tree: connected graph without loops.
- spanning tree: tree containing all the nodes.



Example: Finding the shortest path

- <u>Task</u>: Find the shortest path in graph G(N, A) from s (source) to t (target).
- Edges are labeled with weights:

$$d_{i,j} \in R \Leftrightarrow (i,j) \in A$$

$$d_{i,j} = \infty \Leftrightarrow (i,j) \notin A$$

$$d_{i,i} = 0$$

- Find the *L* path between *s* and *t* where the sum of weights is minimal (shortest path).
 - One possible solution: Dijkstra algorithm



Example: Dijkstra algorithm

Notation:

- D_i = iterated distance of node *i* from s (s=,,1")
- P = set of considered nodes

Algorithm:

- 1. <u>initialization</u>: $P = \{1\}$; $D_j = d_{1j}$ for all j
- 2. <u>iteration</u>: find *i* where $D_i = \min_{\forall j \notin P} D_j$
- 3. let: $P = P \cup \{i\}$ $D_j = \min \{D_j, d_{ij} + D_i\}$
- 4. stop condition: If N = P, STOP else GOTO 2.

